

Secular Equations - Applications and Methods

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with contributions from

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... and **Gene Golub**

$$(A + \lambda B)x(\lambda) = b, \quad \theta(x(\lambda)) = \tau(\lambda)$$

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Summary of the talk

- 🌈 motivation and examples
- 🌈 solving the explicit problem
- 🌈 solving the implicit problem by factorization
- 🌈 solving the implicit problem by iteration
- 🌈 comments and conclusions

History

- **Secular** (adj.) 3. astronomy, of or denoting slow changes in the motion of the sun or planets (c.f. 1. not connected with religious or spiritual matters) (O.E.D.)
- early use of the term **secular equation** in papers by Cauchy and Sylvester refer to equations from which the eigenvalues of a real symmetric matrix may be obtained
- Hilbert (1924) and E.T. Browne (1930) explicitly refer to the characteristic equation as a secular equation
- Golub (SIAM Review, 1973) gives a number of examples of more general secular equations (to follow). See also Golub & Meurant (Matrices, Moments and Quadrature with Applications, 2010)

Secular equations

- given real symmetric A and B and real vector b
- let $x(\lambda)$ be a solution of the linear system

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for given real scalar λ

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- find values of λ that solve the **secular equation**

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for given real **secular functions** $\theta : \mathbb{R}^n \rightarrow \mathbb{R}$ and $\tau : \mathbb{R} \rightarrow \mathbb{R}$

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- often impose extra conditions that identify particular λ

Examples

- eigenvalues of a symmetric bordered matrix
- eigenvalues following a symmetric rank-one change
- the trust-region subproblem
- regularization of quadratic minimization
- global optimization
- total least squares

Eigenvalues of a symmetric bordered matrix

- given real symmetric A , vector b and scalar β , find the eigenvalues of

$$\begin{pmatrix} A & b \\ b^T & \beta \end{pmatrix}$$

- require eigenvalues λ for which

$$\begin{pmatrix} A & b \\ b^T & \beta \end{pmatrix} \begin{pmatrix} v \\ \xi \end{pmatrix} = \lambda \begin{pmatrix} v \\ \xi \end{pmatrix}$$

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- if $\xi = 0 \implies Av = \lambda v$
- otherwise $x = -v/\xi \implies$

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- secular functions

$$\theta(x) = b^T x \text{ and } \tau(\lambda) = \beta - \lambda$$

Eigenvalues following a symmetric rank-one change

- given real symmetric A , vector b and scalar β , find the eigenvalues of

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- prime example: Divide & Conquer method for tridiagonal eigenproblem

(Cuppen, 1981)

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The trust-region subproblem

- given real symmetric H , vector g & radius $\Delta > 0$

$$\underset{x \in \mathbb{R}^n}{\text{minimize}} \quad q(x) = \frac{1}{2}x^T H x + g^T x \quad \text{subject to} \quad \|x\|_2 \leq \Delta$$

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(Gay, Moré, Sorensen, 1981-3)

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- if $\|x_*\|_2 = \Delta$, need to find largest root of $(*)$ with secular functions

$$\theta(x) = \|x\|_2 \quad \text{and} \quad \tau(\lambda) = \Delta$$

Regularization of quadratic minimization

- given real symmetric H , vector g weight $\sigma > 0$ & index $p \geq 2$

$$\underset{x \in \mathbb{R}^n}{\text{minimize}} \quad r(x) = \frac{1}{2}x^T H x + g^T x + \frac{1}{p}\sigma \|x\|_2^p$$

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- need to find largest root of $(*)$ with secular functions

$$\theta(x) = \|x(\lambda)\|_2 \text{ and } \tau(\lambda) = (\lambda/\sigma)^{1/(p-2)}$$

Global Lipschitz optimization

- given a C^2 function $f(x)$ with Lipschitz Hessian, Taylor \implies

$$c^L(s) \equiv f(x) + s^T g(x) + \frac{1}{2} s^T H(x) s - \frac{1}{3} \sigma \|s\|_2^3 \leq \\ f(x + s) \leq f(x) + s^T g(x) + \frac{1}{2} s^T H(x) s + \frac{1}{3} \sigma \|s\|_2^3 \equiv c^U(s)$$

where g/H are the gradient/Hessian of f & σ is the Lipschitz constant

- global minima of c^L and c^U over $\mathcal{S} = \{s : \|s\|_2 \leq \Delta\}$ provide lower and upper bounds for global minimum of $f(x + s)$ over \mathcal{S}
- partition “space of interest” for f into overlapping hyper-spheres and apply branch and bound to find global minimizer in this space
- global minima of c^L and c^U over \mathcal{S} each has associated secular equation \implies tractable

(Cartis, Farmer, Fowkes, G., 2012)

Simplification of secular equations

- without loss of generality consider

$$(A \pm \lambda I)x(\lambda) = b \text{ and } \theta(x(\lambda)) = \tau(\lambda)$$

for real, symmetric A , where for simplicity

$$\theta(x) = b^T x \text{ or } \theta(x) = \|x\|_2$$

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- suppose that eigen-decomposition $A = Q\Lambda Q^T$ is known, with $\Lambda = \text{Diag}(\lambda_i)$: $\lambda_i \leq \lambda_{i+1}$ & orthogonal Q , and that $\bar{b} = Q^T b \implies x(\lambda) = Q\bar{x}(\lambda)$:

$$(\Lambda \pm \lambda I)\bar{x}(\lambda) = \bar{b} \text{ and } \bar{\theta}(\bar{x}(\lambda)) = \tau(\lambda)$$

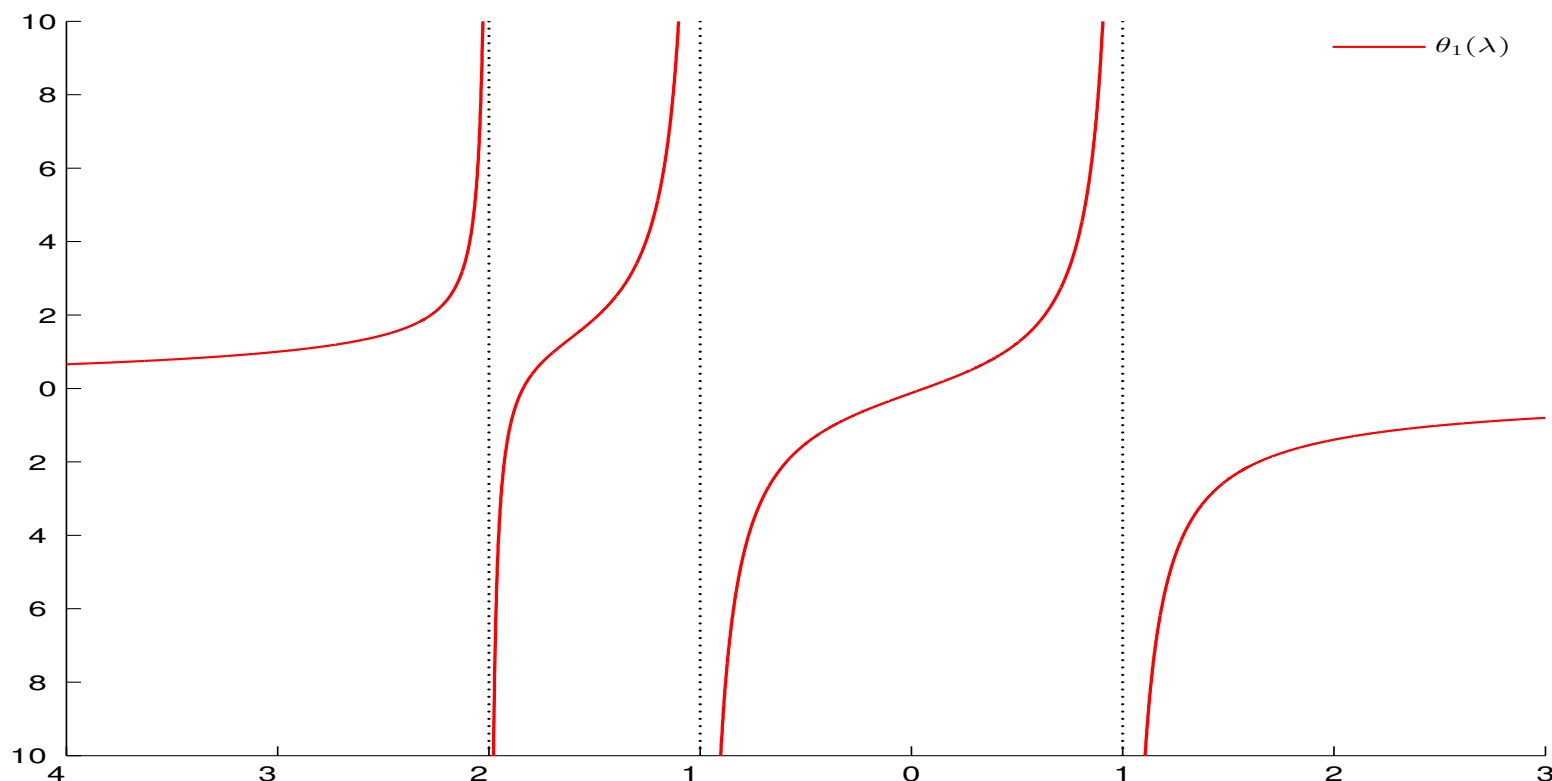
and, e.g.,

$$\bar{\theta}(\bar{x}) = \bar{b}^T \bar{x} \text{ or } \bar{\theta}(\bar{x}) = \|\bar{x}\|_2$$

Type-I secular equations

$$(\Lambda - \lambda I)\bar{x}(\lambda) = \bar{b} \text{ and } \theta_1(\lambda) \equiv \bar{b}^T \bar{x}(\lambda) = \tau(\lambda) \implies$$

$$\theta_1(\lambda) = \sum_{i=1}^n \frac{\bar{b}_i^2}{\lambda_i - \lambda}$$

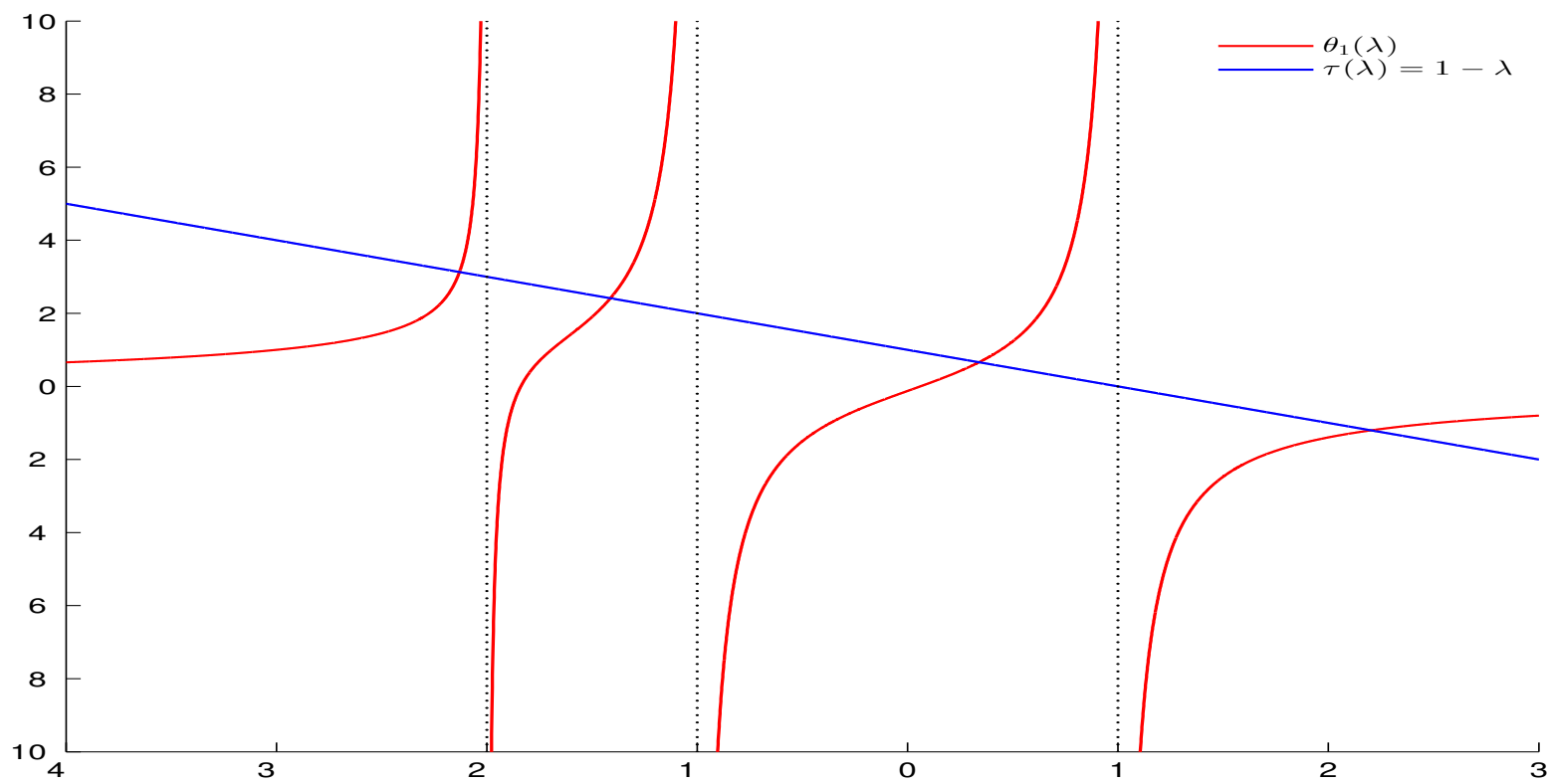


The type-I secular function $\theta_1(\lambda)$

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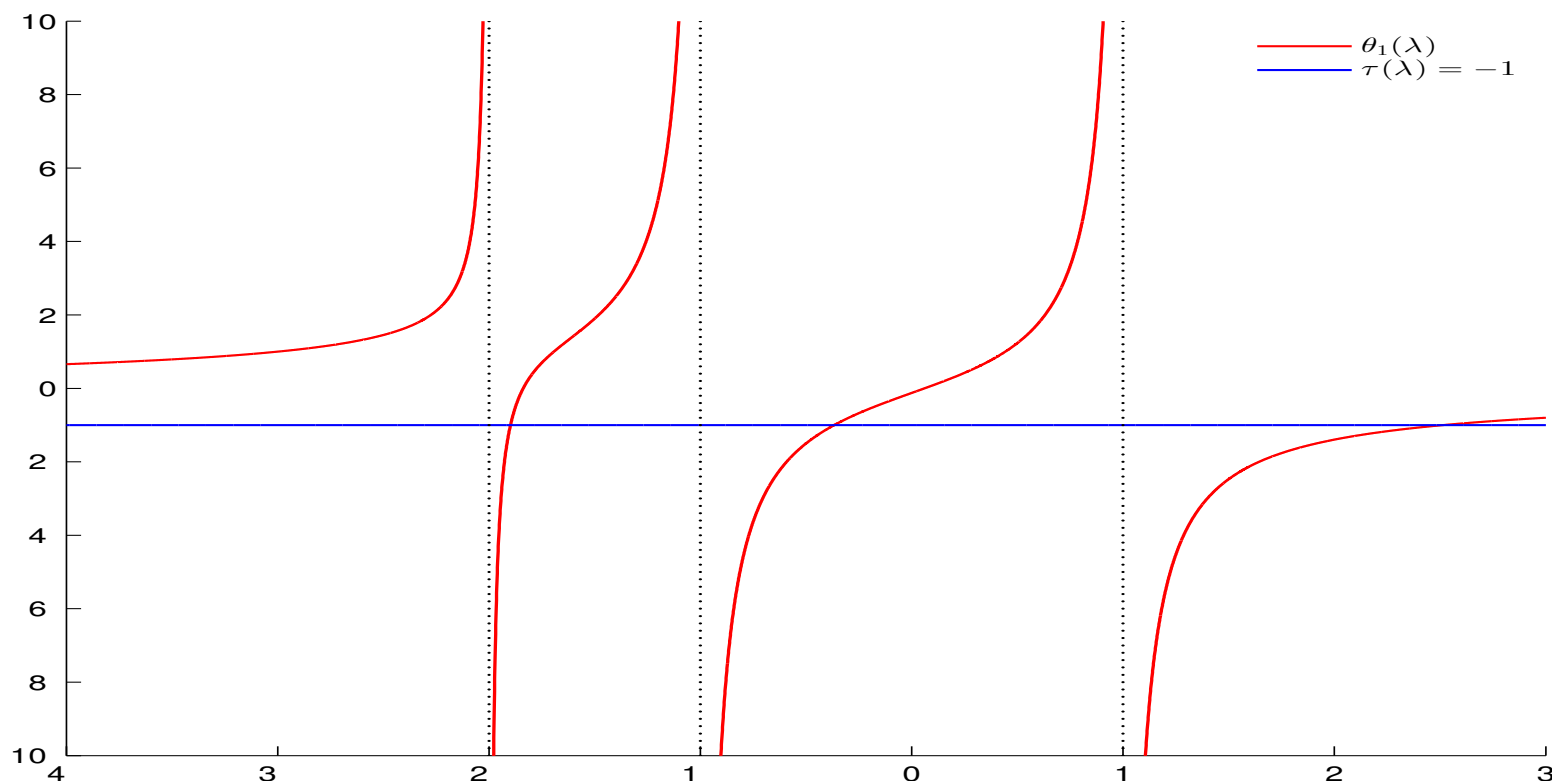


The type-I secular equation for the eigenvalues of a symmetric bordered matrix

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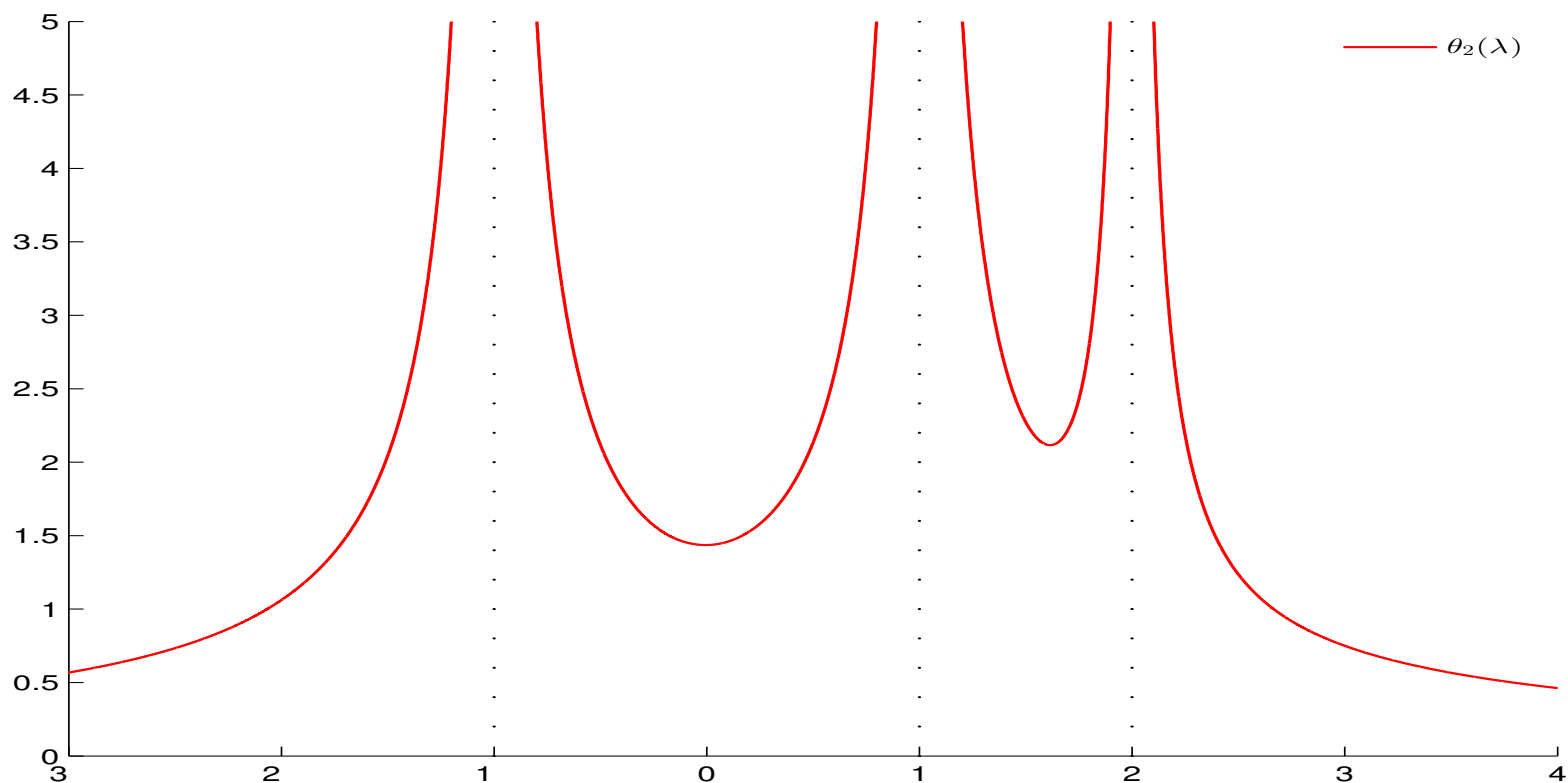


The type-I secular equation for the eigenvalues following a symmetric rank-one change

Type-II secular equations

$$(\Lambda + \lambda I)\bar{x}(\lambda) = \bar{b} \text{ and } \theta_2(\lambda) \equiv \|\bar{x}(\lambda)\|_2 = \tau(\lambda) \implies$$

$$\theta_2(\lambda) = \sqrt{\sum_{i=1}^n \left(\frac{\bar{b}_i}{\lambda_i + \lambda} \right)^2}$$

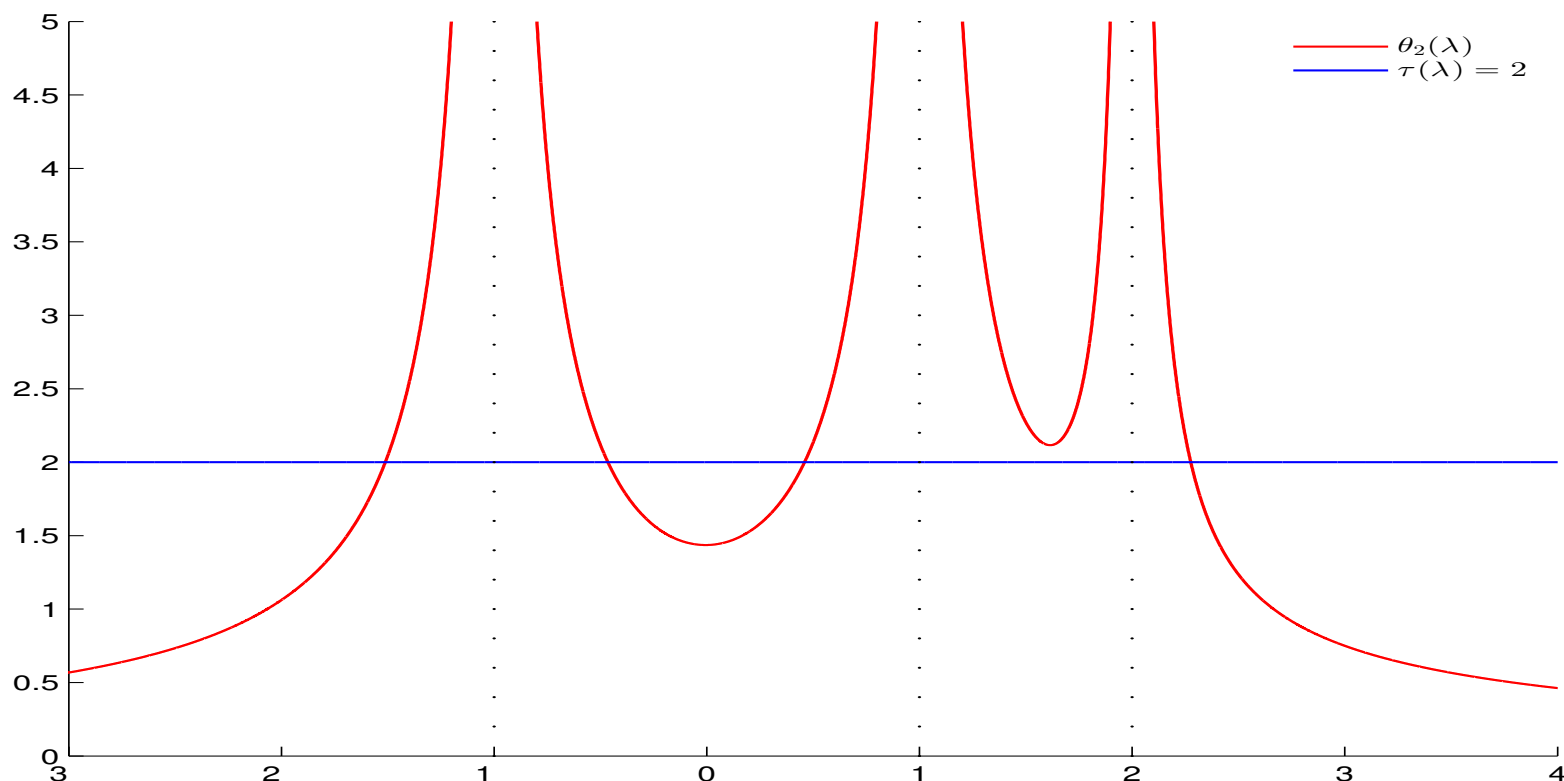


The type-II secular function $\theta_2(\lambda)$

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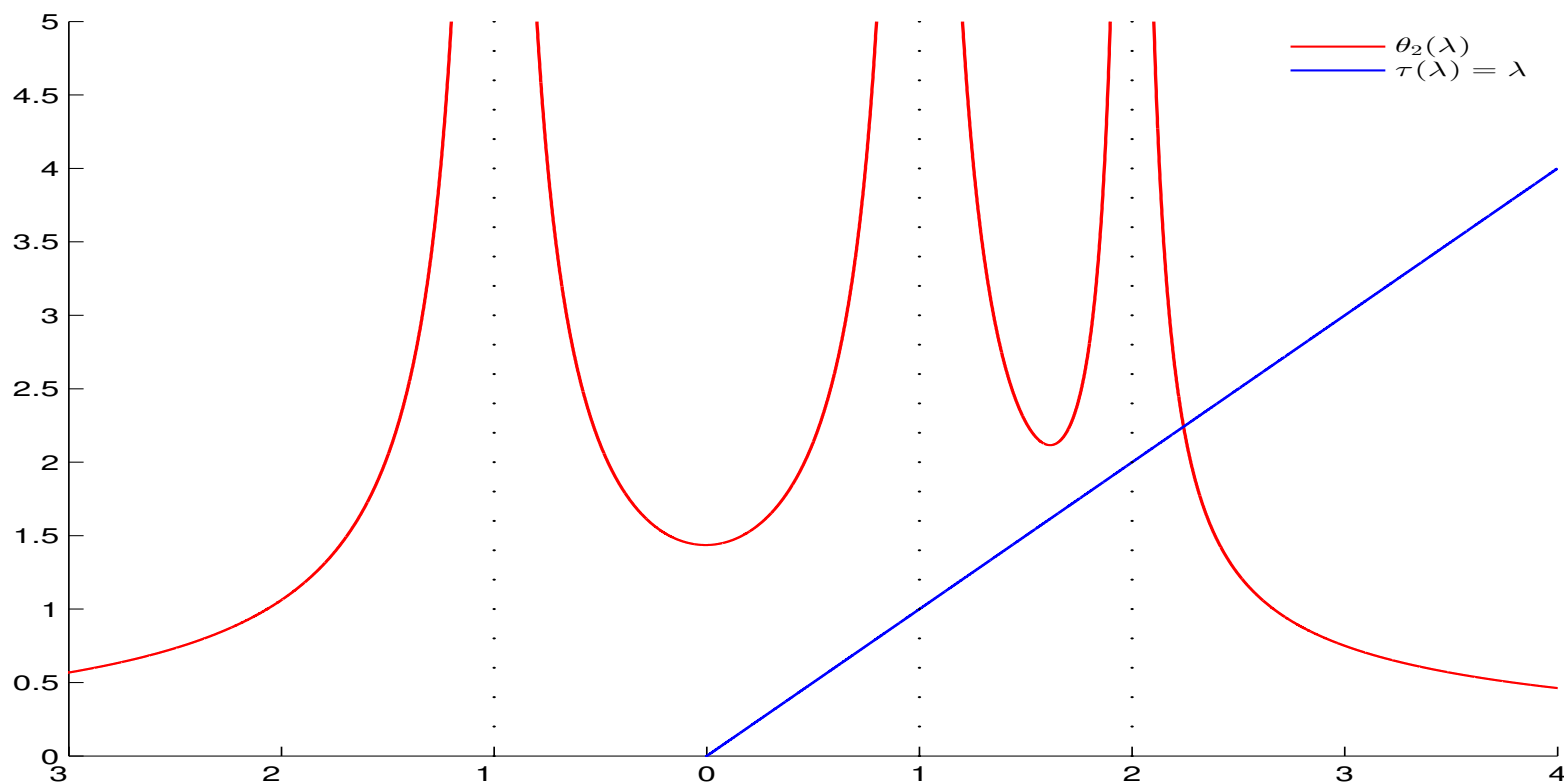


The type-II secular equation for the trust-region subproblem

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The type-II secular equation for cubic regularization of quadratic minimization

Explicit vs Implicit

In general most secular equations occur as one of two types:

🌈 **Type I** secular equation:

$$(A - \lambda I)x(\lambda) = b \text{ and } \theta_1(\lambda) \equiv b^T x(\lambda) = \tau(\lambda)$$

- 🌈 usually more than one—maybe all—roots required
- 🌈 often eigenvalues/vectors known \implies reduces to explicit solution of

$$\sum_{i=1}^n \frac{d_i}{\lambda_i - \lambda} = \tau(\lambda) \text{ with } d_i > 0$$

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• **Type II** secular equation:

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$\tau(\lambda)$ usually a low-order (constant or linear) polynomial

Solving the explicit Type-I problem

Possible ways to find root $\in (\lambda_j, \lambda_{j+1})$ of

$$p_1(\lambda) + p_2(\lambda) \equiv \sum_{i=1}^j \frac{d_i}{\lambda_i - \lambda} + \sum_{i=j+1}^n \frac{d_i}{\lambda_i - \lambda} = \tau(\lambda) \quad \text{with } d_i > 0$$

- construct low-order rational approximations $r_1^{(k)}$ and $r_2^{(k)}$ to p_1 and p_2 with poles λ_j and λ_{j+1}
- find appropriate root $\lambda^{(k)}$ of low-order polynomial equation

$$r_1^{(k)}(\lambda) + r_2^{(k)}(\lambda) = \tau(\lambda)$$

- appropriate approximations lead to global quadratic convergence of $\lambda^{(k)}$ from left or right of root, or locally from either side

(Bunch, Nielsen, Sorensen, 1978, Li, 1994)

Solving the explicit Type-I problem (continued)

Possible ways to find root $\in (\lambda_j, \lambda_{j+1})$ of

$$\sum_{i=1}^n \frac{d_i}{\lambda_i - \lambda} = \tau(\lambda) \quad \text{with } d_i > 0$$

- use nonlinear change of variable $\lambda = \lambda_j + 1/\mu$
- apply standard fast method (Newton, Gragg, Halley, ...) to problem in transformed variable (better-behaved problem)
- superlinear to cubic global convergence established (Melman, 1995-97)

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- accumulated evidence of good performance in practice (3-5 iterations for high accuracy)

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- accumulated evidence of good performance in practice (3-5 iterations for high accuracy)
- $O(n)$ approximation methods for **all** roots related to multi-level integral transform evaluation (Livne, Brandt, 2002)

Solving the implicit Type-II problem

Find the largest root $\geq \max(0, -\lambda_1)$ of

$$\theta_2(\lambda) \equiv \|x(\lambda)\|_2 = \tau(\lambda) \text{ where } (A + \lambda I)x(\lambda) = b$$

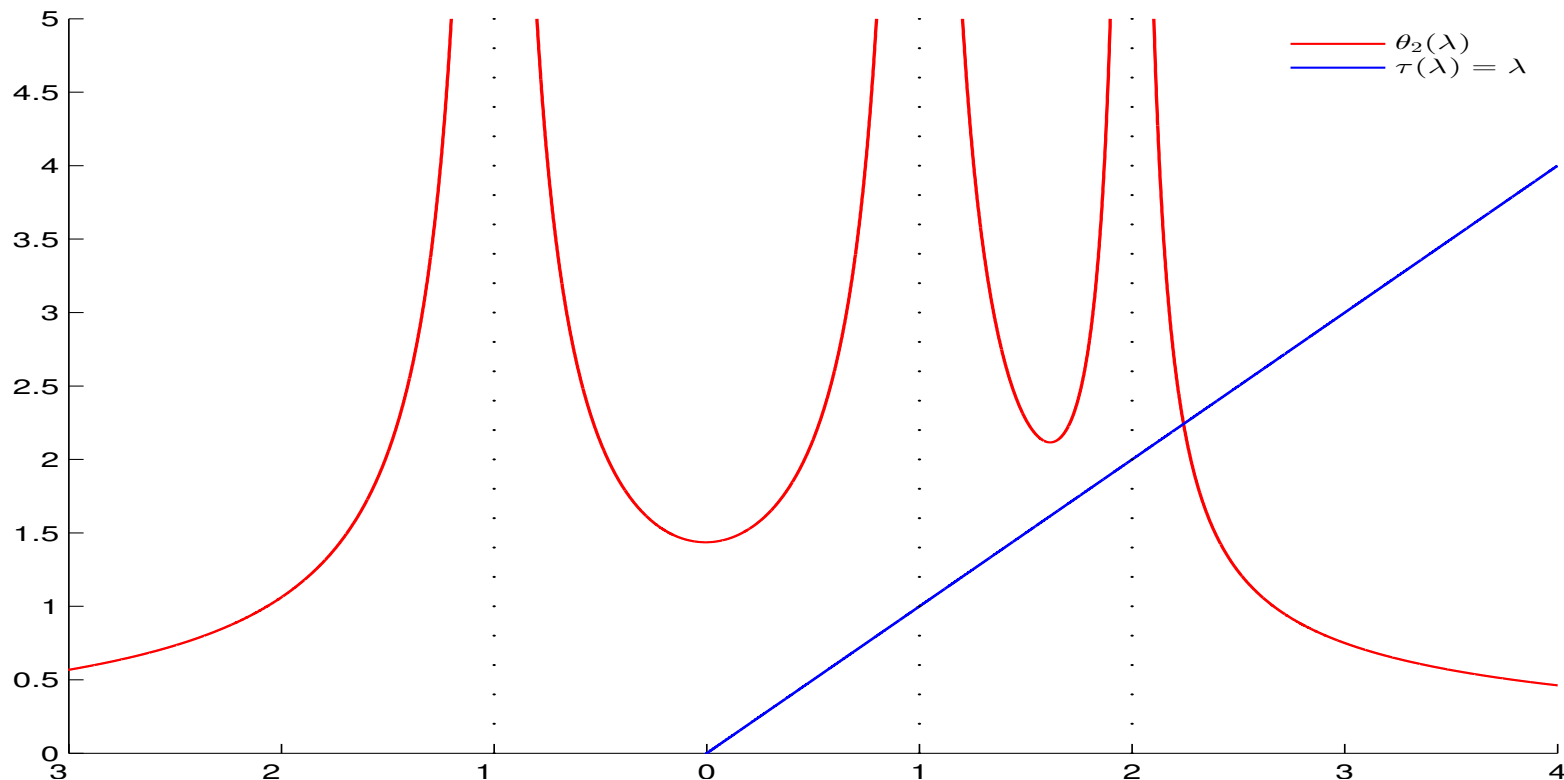
- usually $\tau(\lambda)$ is convex, increasing for positive λ
 \implies unique root in $[\max(0, -\lambda_1), \infty)$, but generally λ_1 unknown
- need to use effective rootfinder that “understands” structure of $\theta_2(\lambda)$
- need to evaluate $\theta_2(\lambda)$ and its derivatives

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🌈 “degenerate” **hard** case if $b \perp$ eigenvector(s) q_1 for λ_1



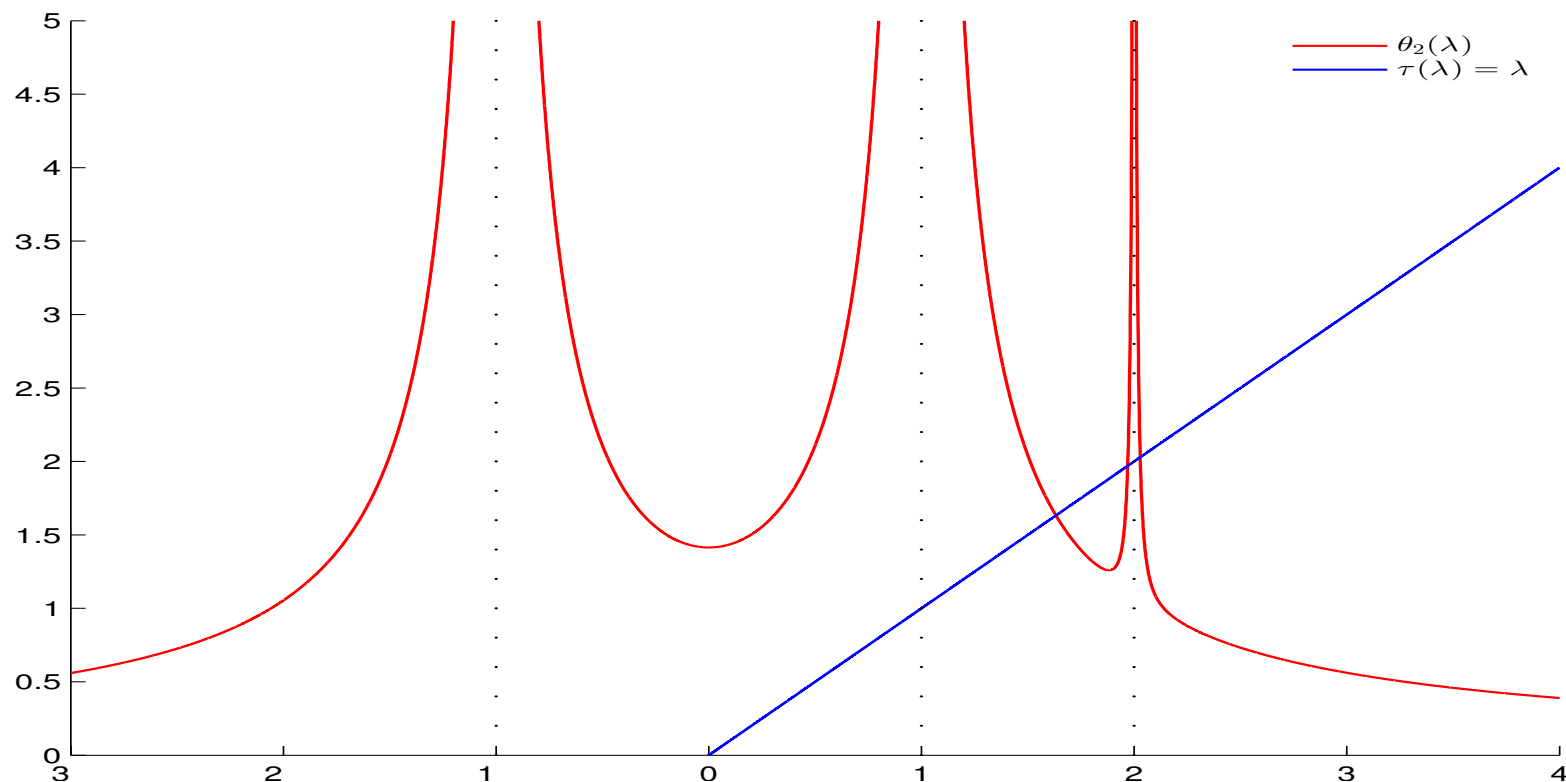
The type-II secular equation for cubic regularization of quadratic minimization - easy case

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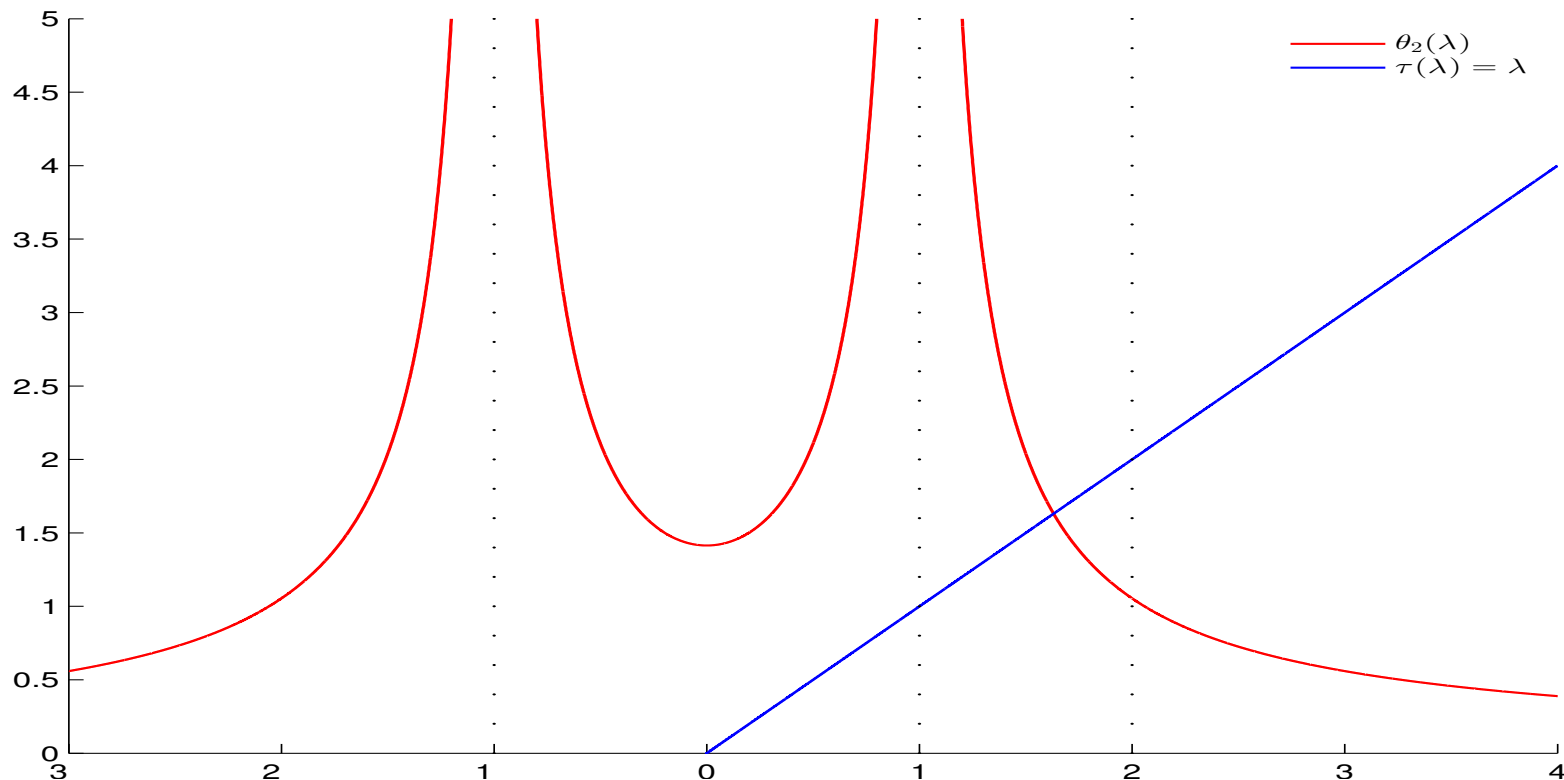
The type-II secular equation for cubic regularization of quadratic minimization - moderate case

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The type-II secular equation for cubic regularization of quadratic minimization - hard case

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- “degenerate” **hard** case if $b \perp$ eigenvector(s) q_1 for λ_1
- solution in hard case is $x(-\lambda_1) + \alpha q_1$ where α satisfies

$$\|x(-\lambda_1) + \alpha q_1\|_2 = \tau(-\lambda_1)$$

Solving the implicit Type-II problem

Find the largest root $\geq \max(0, -\lambda_1)$ of

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🌈 may equivalently find the largest root $\geq \max(0, -\lambda_1)$ of

$$\pi(\lambda) \equiv \|x(\lambda)\|_2^2 = \tau^2(\lambda) \equiv \rho(\lambda) \text{ where } (A + \lambda I)x(\lambda) = b$$

since derivatives of $\pi(\lambda)$ are easy to find

Calculating the secular function and its derivatives

Need to compute

$$\pi(\lambda) \equiv \|x(\lambda)\|_2^2 = x^T(\lambda)x(\lambda) \text{ where } (A + \lambda I)x(\lambda) = b$$

and its derivatives

🌈 derivatives $x^{(k)}(\lambda)$ of $x(\lambda)$

$$(A + \lambda I)x^{(1)} = -x, \quad (A + \lambda I)x^{(2)} = -2x^{(1)}$$

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$$\begin{aligned}(A + \lambda I)x^{(1)} &= -x, & (A + \lambda I)x^{(2)} &= -2x^{(1)} \\ \dots & & (A + \lambda I)x^{(k)} &= -kx^{(k-1)}\end{aligned}$$

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$$\pi(\lambda) \equiv \|x(\lambda)\|_2^2 = x^T(\lambda)x(\lambda) \text{ where } (A + \lambda I)x(\lambda) = b$$

and its derivatives

🌈 derivatives $x^{(k)}(\lambda)$ of $x(\lambda)$

$$(A + \lambda I)x^{(1)} = -x, \quad (A + \lambda I)x^{(2)} = -2x^{(1)}$$

🌈 derivatives $\pi^{(k)}(\lambda)$ of $\pi(\lambda)$

$$\pi^{(1)} = 2x^{(1)T}x, \quad \pi^{(2)} = 2x^{(2)T}x + 2x^{(1)T}x^{(1)} = 6x^{(1)T}x^{(1)}$$

Calculating the secular function and its derivatives

Need to compute

$$\pi(\lambda) \equiv \|x(\lambda)\|_2^2 = x^T(\lambda)x(\lambda) \text{ where } (A + \lambda I)x(\lambda) = b$$

and its derivatives

🌈 derivatives $x^{(k)}(\lambda)$ of $x(\lambda)$

$$(A + \lambda I)x^{(k)} = -kx^{(k-1)}$$

🌈 derivatives $\pi^{(k)}(\lambda)$ of $\pi(\lambda)$

$$\pi^{(1)} = 2x^{(1)T}x, \quad \pi^{(2)} = 6x^{(1)T}x^{(1)}$$

$$\dots \pi^{(2k+1)} = 2\alpha_k x^{(k)T}x^{(k+1)}, \quad \pi^{(2k+2)} = \alpha_{k+1} x^{(k+1)T}x^{(k+1)}$$

$$\text{where } \alpha_{k+1} = 2(2k+3)\alpha_k/(k+1)$$

Calculating the secular function and its derivatives

Need to compute

$$\pi(\lambda) \equiv \|x(\lambda)\|_2^2 = x^T(\lambda)x(\lambda) \text{ where } (A + \lambda I)x(\lambda) = b$$

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- derivatives $x^{(k)}(\lambda)$ of $x(\lambda)$

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- derivatives $\pi^{(k)}(\lambda)$ of $\pi(\lambda)$

$$\pi^{(1)} = 2x^{(1)T}x, \quad \pi^{(2)} = 6x^{(1)T}x^{(1)}$$

- given the Cholesky factorization $A + \lambda I = LL^T$ and

$$Ly^{(1)} = -x, \quad L^T x^{(1)} = y^{(1)} \implies \\ \pi^{(1)} = -2\|y^{(1)}\|_2^2, \quad \pi^{(2)} = 6\|x^{(1)}\|_2^2$$

Calculating the secular function and its derivatives

Need to compute

$$\pi(\lambda) \equiv \|x(\lambda)\|_2^2 = x^T(\lambda)x(\lambda) \text{ where } (A + \lambda I)x(\lambda) = b$$

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- derivatives $x^{(k)}(\lambda)$ of $x(\lambda)$

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$$\pi^{(2k+1)} = 2\alpha_k x^{(k)T} x^{(k+1)}, \quad \pi^{(2k+2)} = \alpha_{k+1} x^{(k+1)T} x^{(k+1)}$$

$$\text{where } \alpha_{k+1} = 2(2k + 3)\alpha_k / (k + 1)$$

- given the Cholesky factorization $A + \lambda I = LL^T$ and

$$Ly^{(k)} = -kx^{(k-1)}, \quad L^T x^{(k)} = y^{(k)} \implies$$

$$\pi^{(2k+1)} = -2\beta_k \|y^{(k)}\|_2^2, \quad \pi^{(2k+2)} = \alpha_{k+1} \|x^{(k)}\|_2^2$$

$$\text{where } \beta_k = 2\alpha_k / (k + 1) \text{ and } \alpha_{k+1} = (2k + 3)\beta_k$$

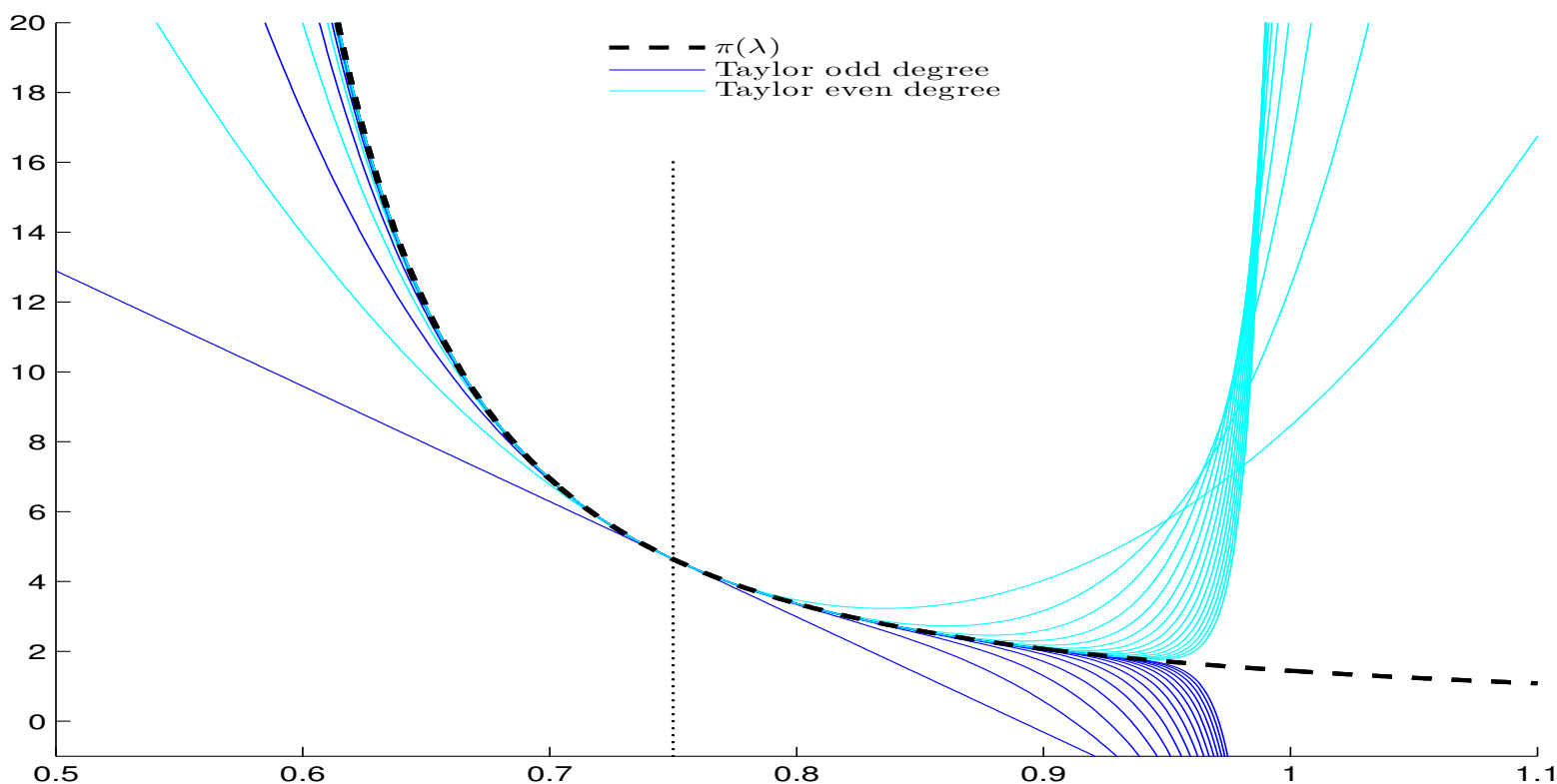
- one triangular solve per order of derivative

Taylor-series approximations

Let $p_i(\lambda)$ be the i th order Taylor approximation to $\pi(\lambda)$ at $\lambda_c > -\lambda_1 \implies$

- $p_{2k+1}(\lambda)$ underestimates $\pi(\lambda)$ for $\lambda > -\lambda_1$ for all k
- $p_k(\lambda)$ underestimates $\pi(\lambda)$ for $-\lambda_1 < \lambda < \lambda_c$ for all k

(G., Robinson, Thorne, 2011)



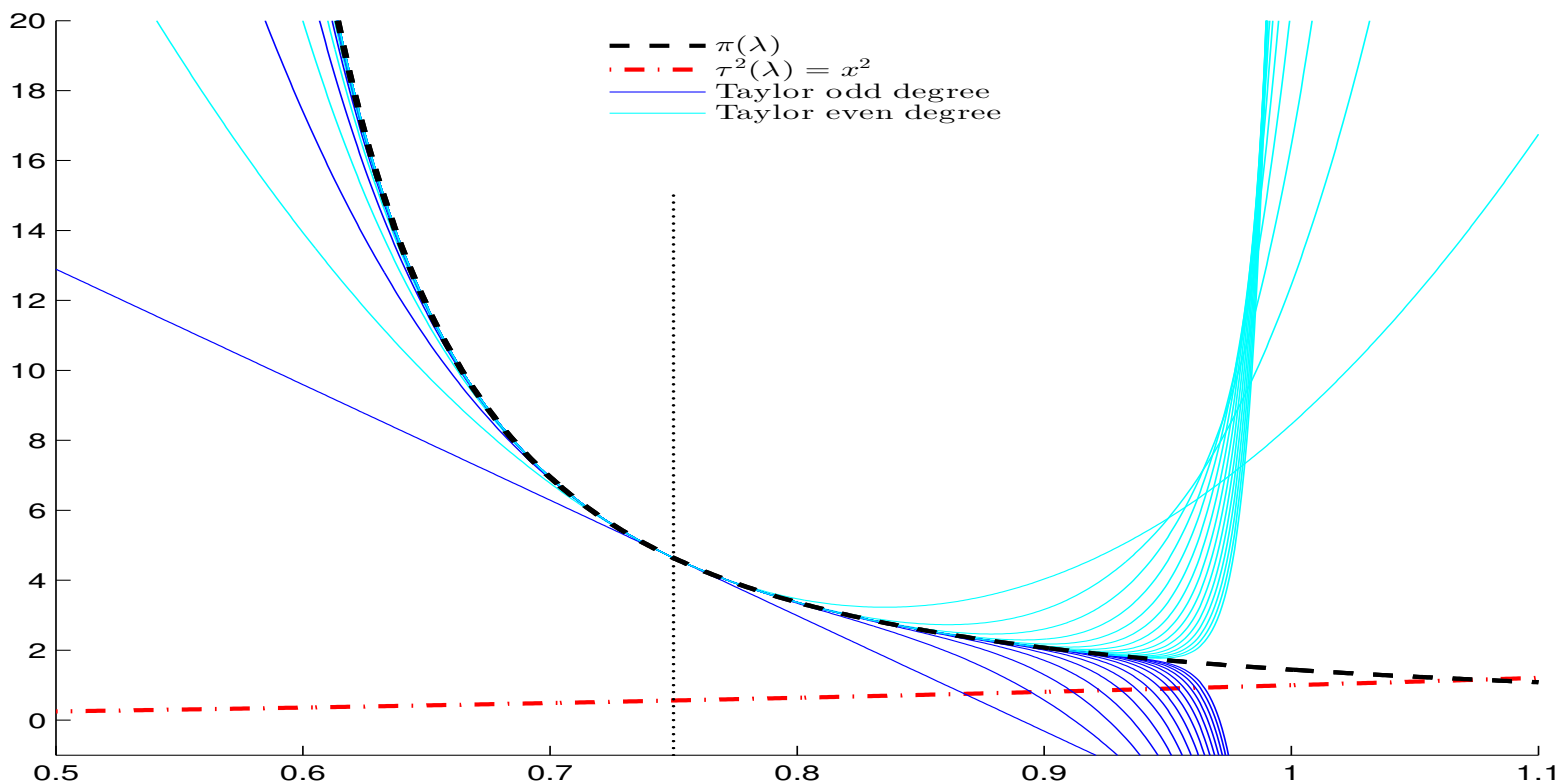
$\pi(\lambda)$ and its first 30 Taylor-series approximations $p_i(\lambda)$ at $\lambda_c = 0.75$

Taylor-series approximations

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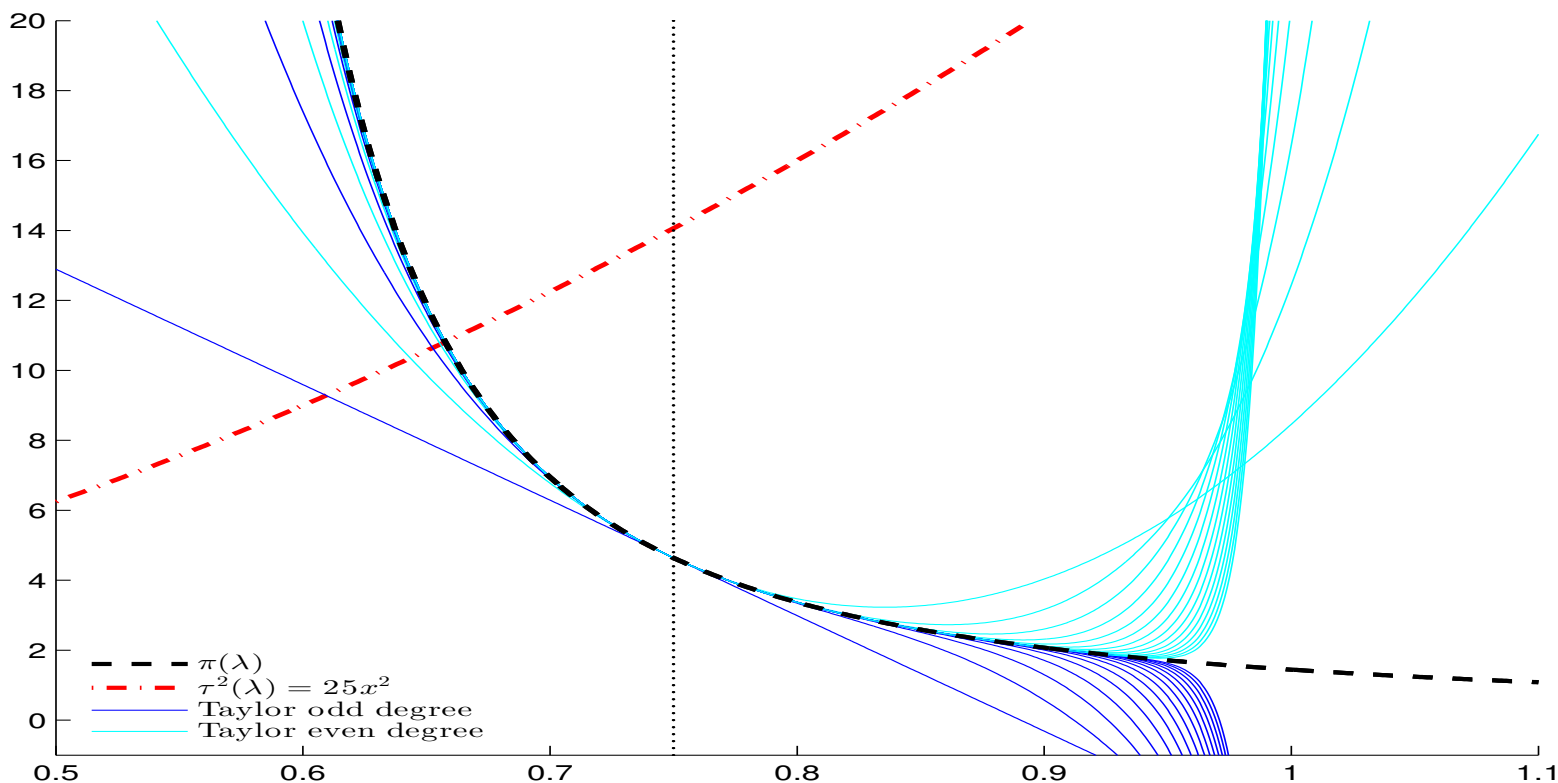
Root of $p_{2k+1}(\lambda) = \rho(\lambda) < \text{root } \lambda_* \text{ of } \pi(\lambda) = \rho(\lambda)$ for $\rho(\lambda) \nearrow$ when $\lambda_c < \lambda_*$

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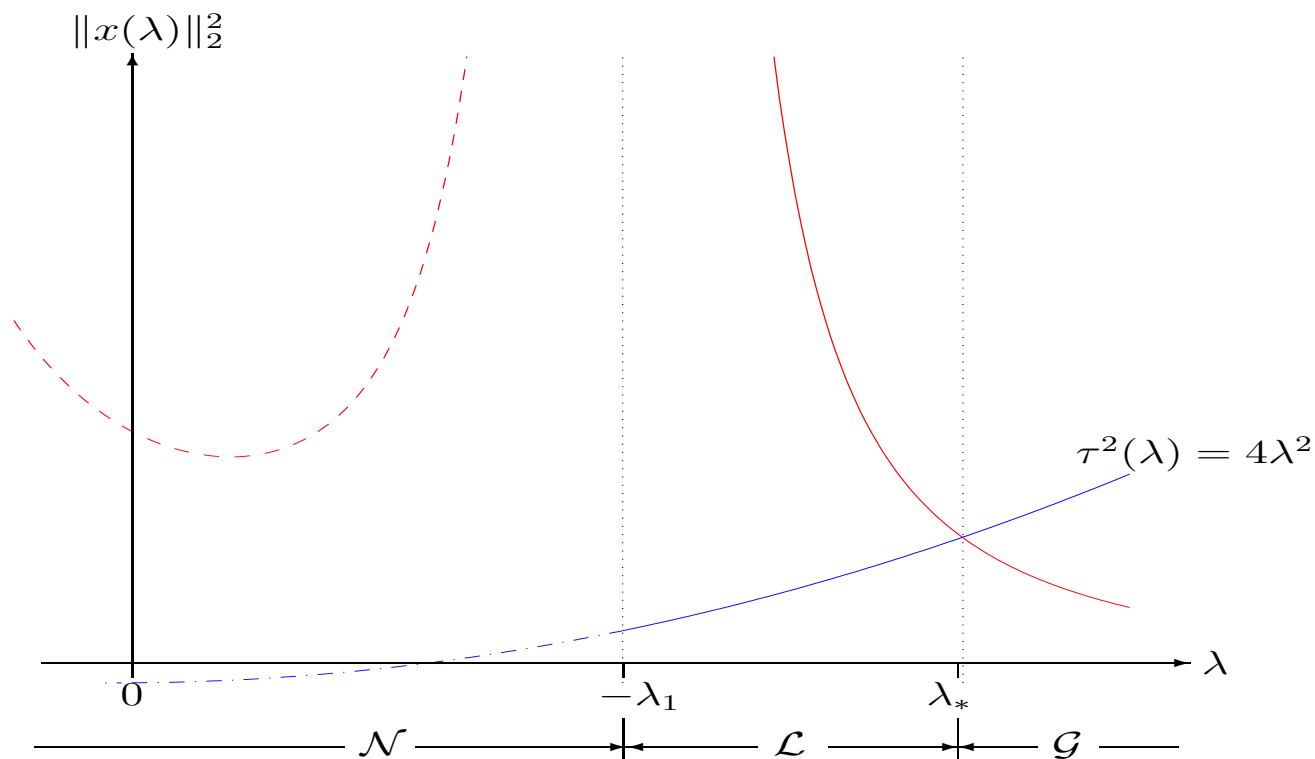
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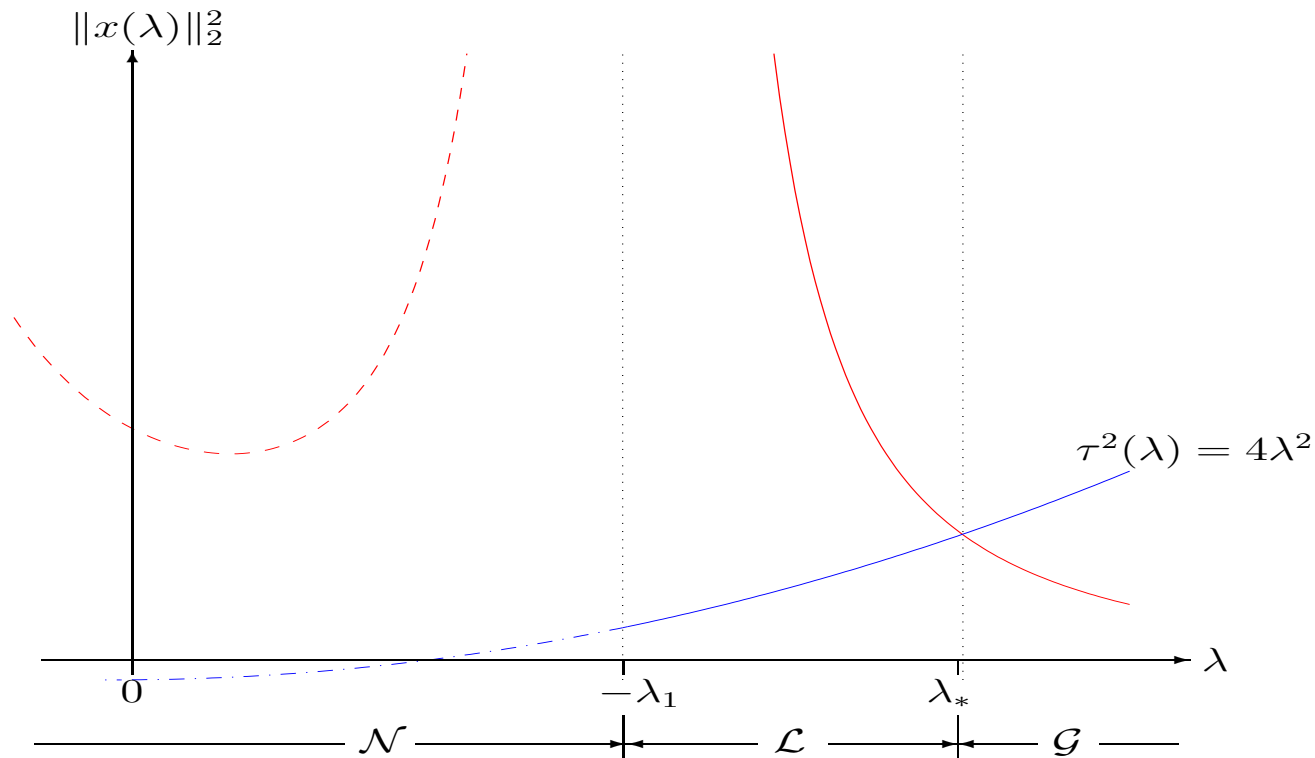
Root of $p_k(\lambda) = \rho(\lambda) < \text{root } \lambda_* \text{ of } \pi(\lambda) = \rho(\lambda)$ for $\rho(\lambda) \nearrow$ when $\lambda_* < \lambda_c$

Algorithm overview



- iterate root of $p_{2k+1}(\lambda) = \tau^2(\lambda)$ within $\mathcal{L} \implies$ globally convergent with Q-order $2k + 2$ (G., Robinson, Thorne, 2011)

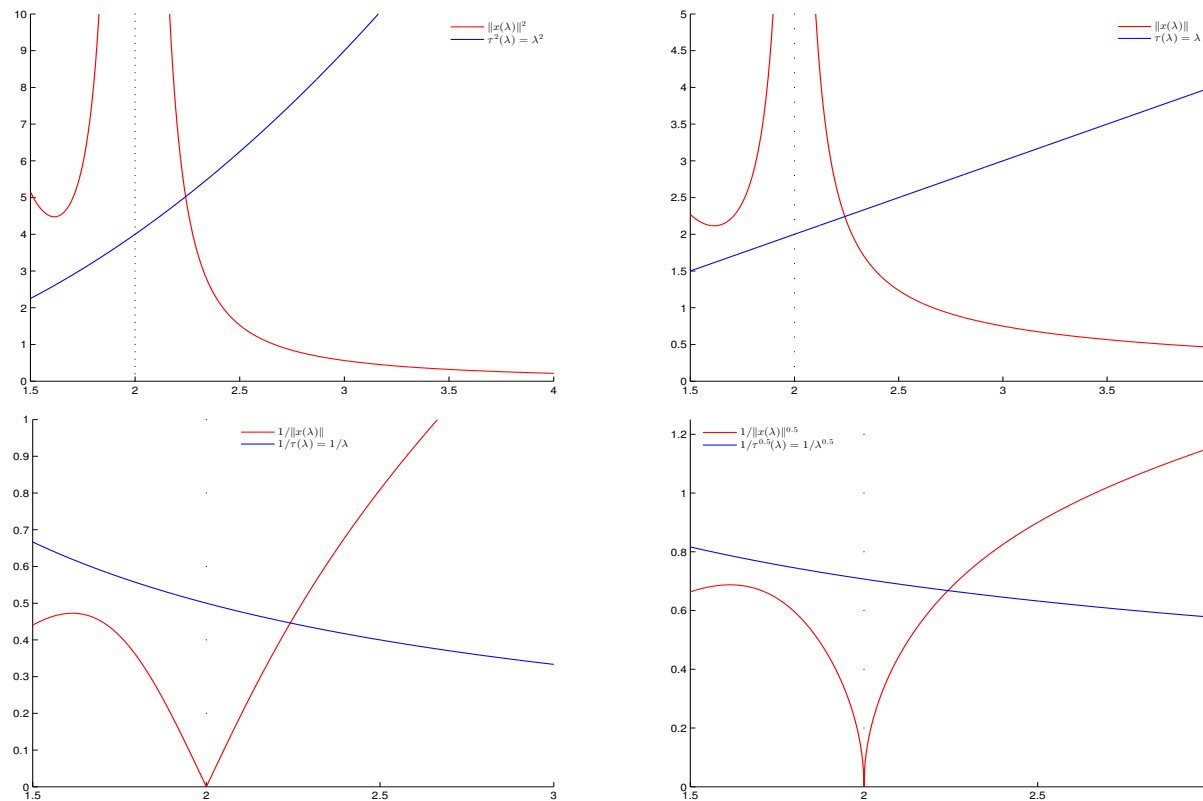
Algorithm overview



- iterate root of $p_{2k+1}(\lambda) = \tau^2(\lambda)$ within $\mathcal{L} \implies$ globally convergent with Q-order $2k + 2$ (G., Robinson, Thorne, 2011)
- safeguarded bisection to locate \mathcal{L}
 - $\lambda \in \mathcal{N} \iff H + \lambda I \neq 0$
 - $\lambda \in \mathcal{G}$ when $\pi(\lambda) < \tau^2(\lambda) \implies$ root of $p_k(\lambda) = \tau^2(\lambda) \notin \mathcal{G}$

Nonlinear transformation

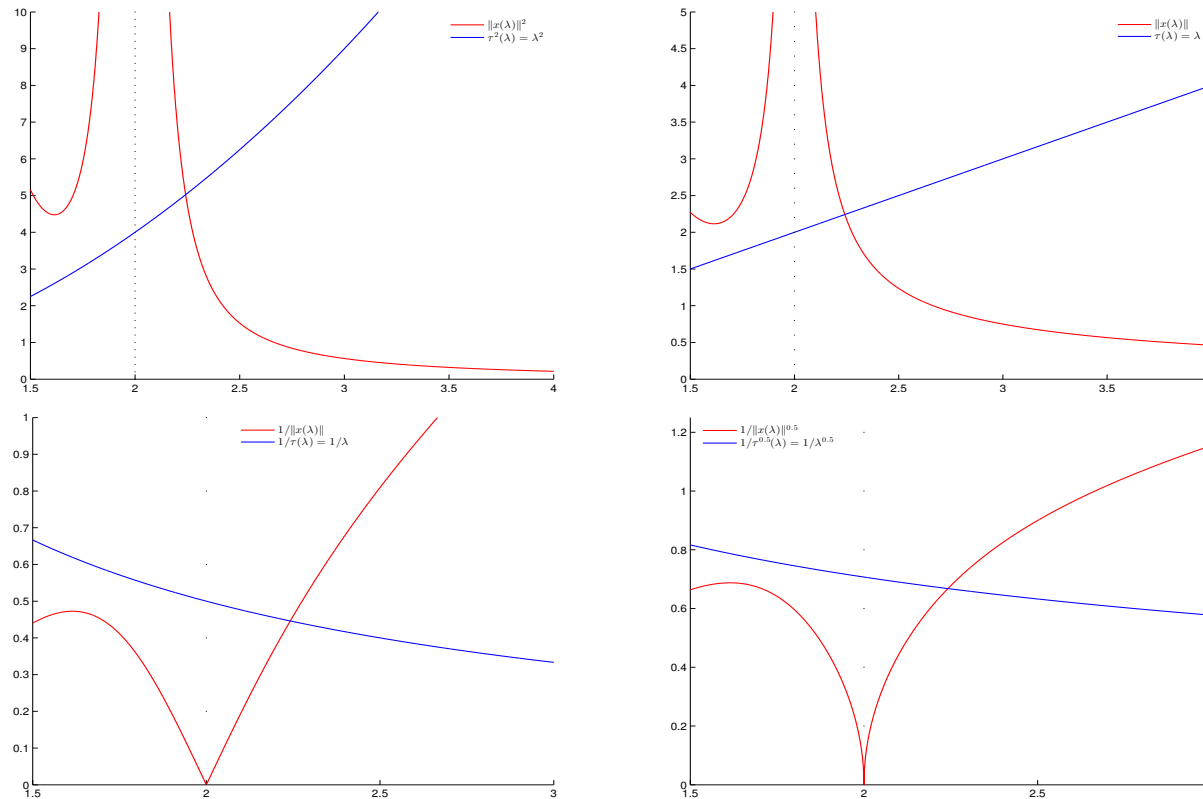
- for $\lambda > -\lambda_1$: $\|x(\lambda)\|_2^\beta$ convex for $\beta > 0$, concave for $\beta \in [-1, 0)$
- solve instead $\|x(\lambda)\|_2^\beta = \tau^\beta(\lambda)$ (Reinsch, 1971)



Clockwise from top left: $\beta = 2, 1, -0.5, -1$

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Clockwise from top left: $\beta = 2, 1, -0.5, -1$

- Taylor approximation $p_1(\lambda)$ to $\|x(\lambda)\|_2^\beta$ for $\beta \geq -1$ underestimates root for $\lambda \in (-\lambda_1, \lambda_*)$ (G., Robinson, Thorne, 2011)

Solving the implicit problem by iteration

What if factorization of $H + \lambda I$ is not possible?

- pick $n \times k$ matrices Q_k with orthonormal columns
- look for approximation to the root $x = Q_k x_k$, where

$$(H + \lambda I)Q_k x_k(\lambda) = b \text{ and } \|Q_k x_k(\lambda)\|_2 = \tau(\lambda)$$

\Rightarrow

$$(Q_k^T H Q_k + \lambda I)x_k(\lambda) = Q_k^T b \text{ and } \|x_k(\lambda)\|_2 = \tau(\lambda)$$

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- if Q_k is generated by the **Lanczos** method (G., Lucidi, Roma, Toint, 1999)
- $Q_k^T H Q_k = \text{tridiagonal } T_k$ and $Q_k^T b = \|b\|_2 e_1$
- $(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \text{ and } \|x_k(\lambda)\|_2 = \tau(\lambda)$

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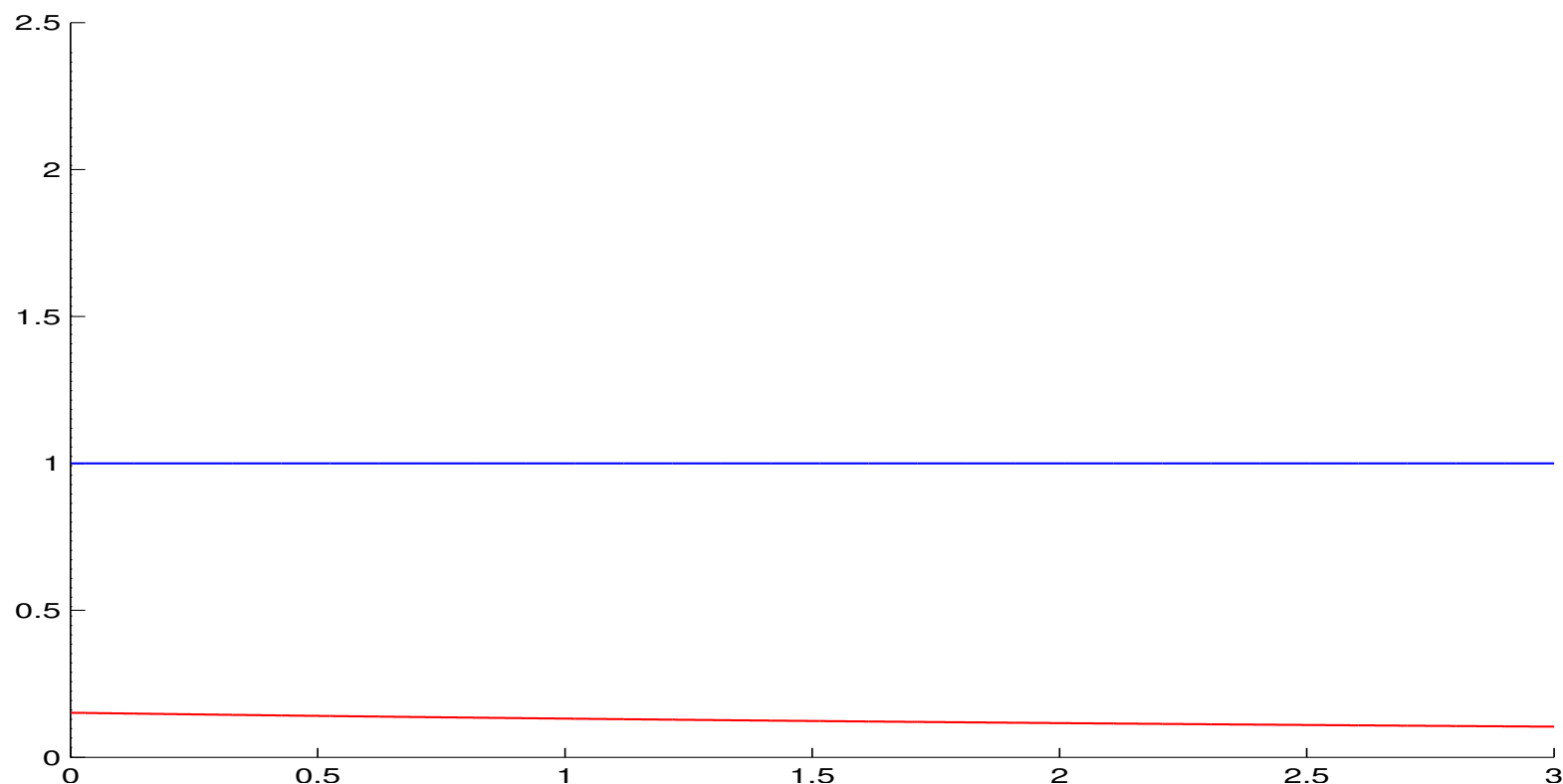
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 - $(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \text{ and } \|x_k(\lambda)\|_2 = \tau(\lambda)$
 - Cholesky factorization of $T_k + \lambda I$ possible in $O(k)$ flops
 - $\|(H + \lambda_k I)Q_k x_k - b\|_2 = \gamma_k e_k^T x_k$, where $\gamma_k = (T_k)_{k,k-1}$
 - good starting guess for λ_{k+1} from k th problem

How the secular equation evolves

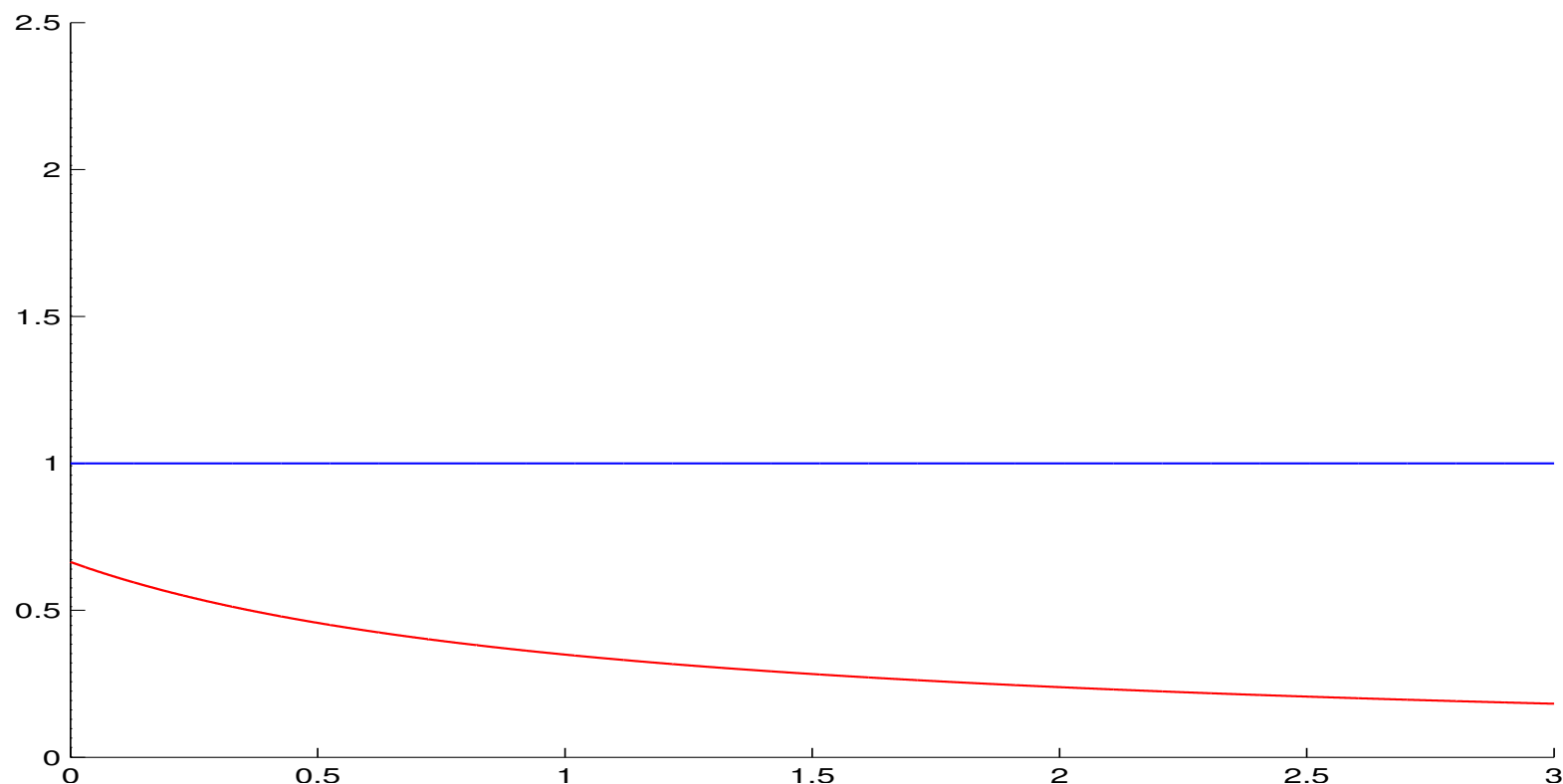
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration k = 1 for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

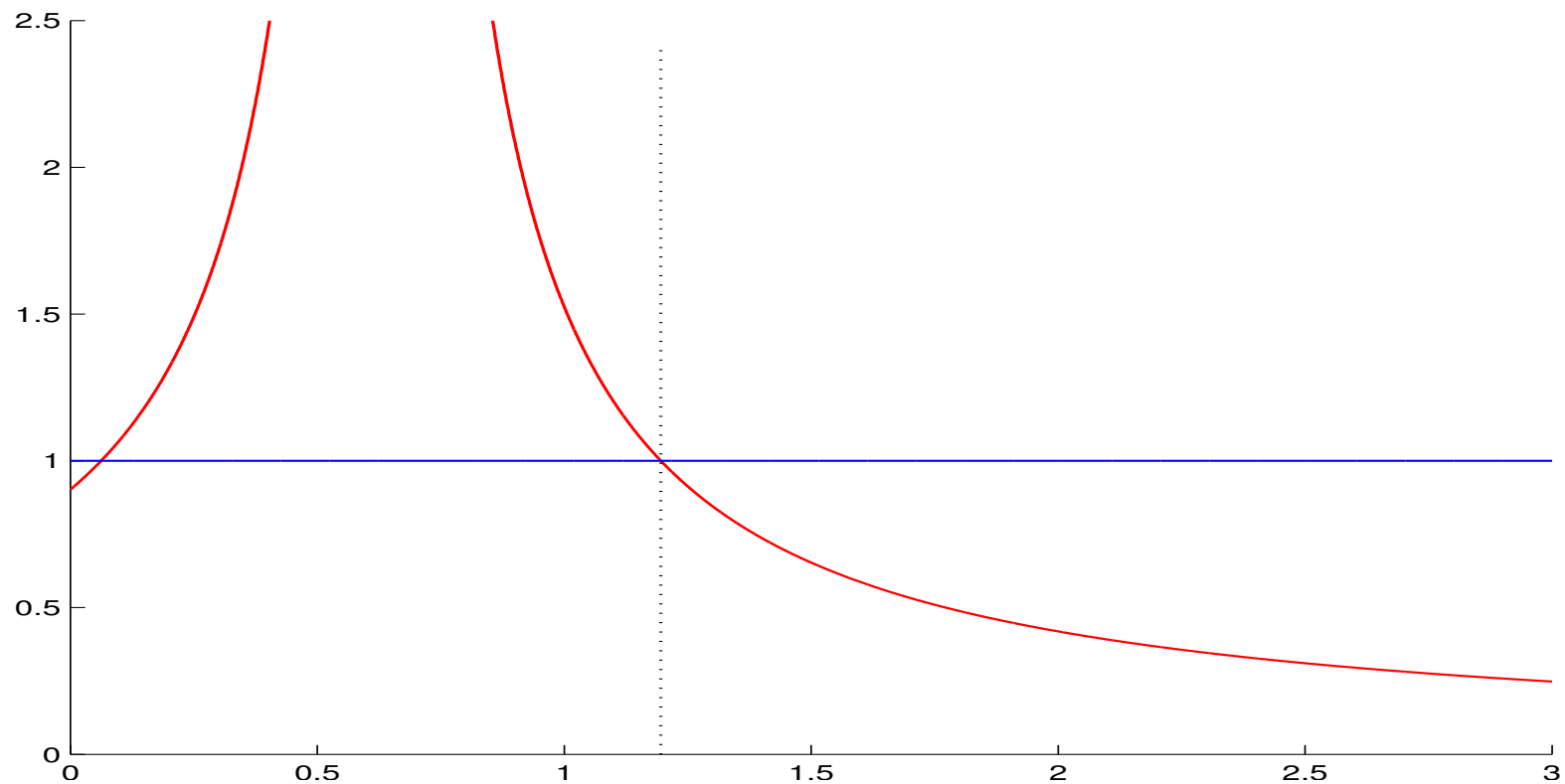
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 2$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

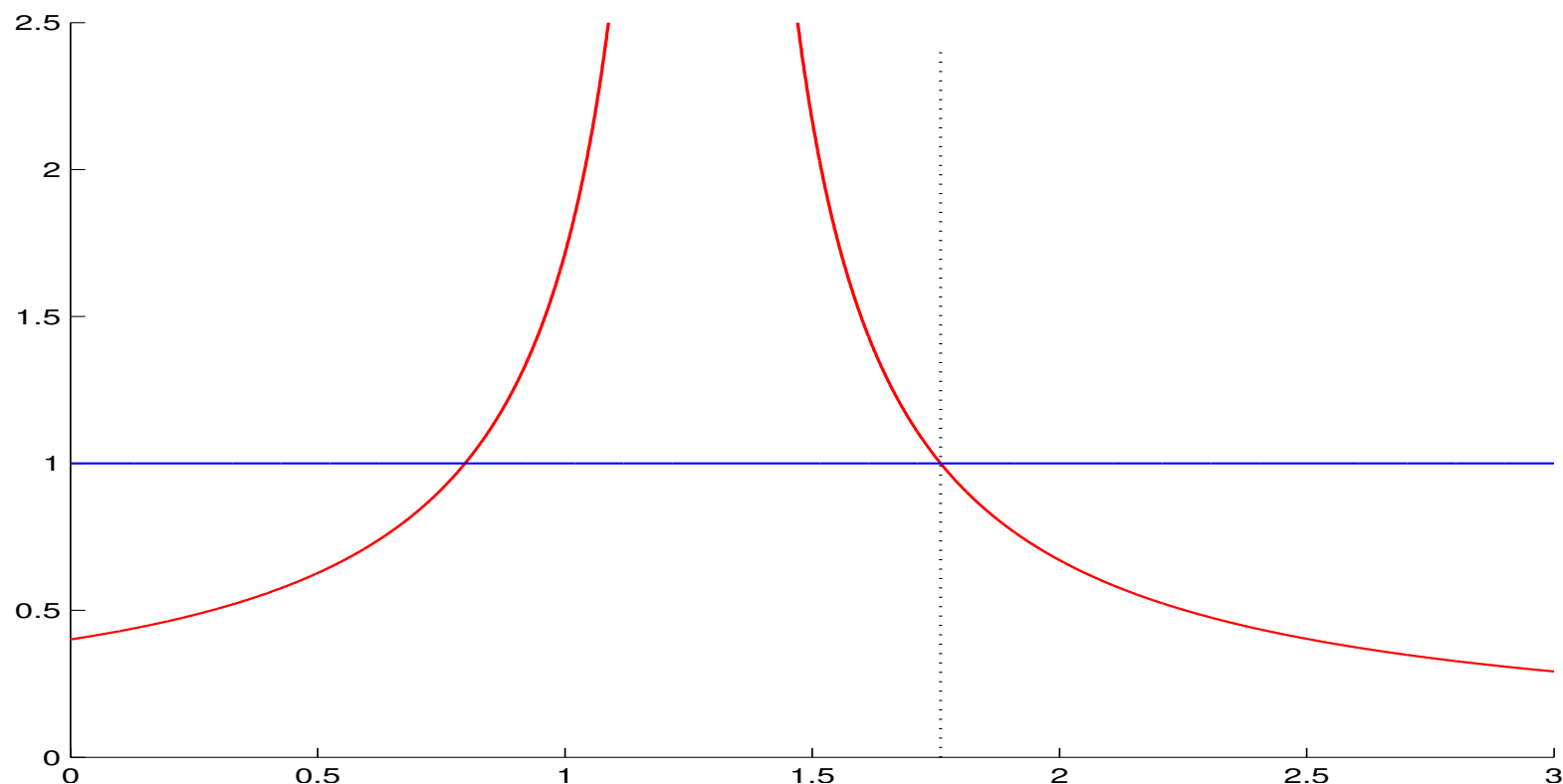
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 3$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

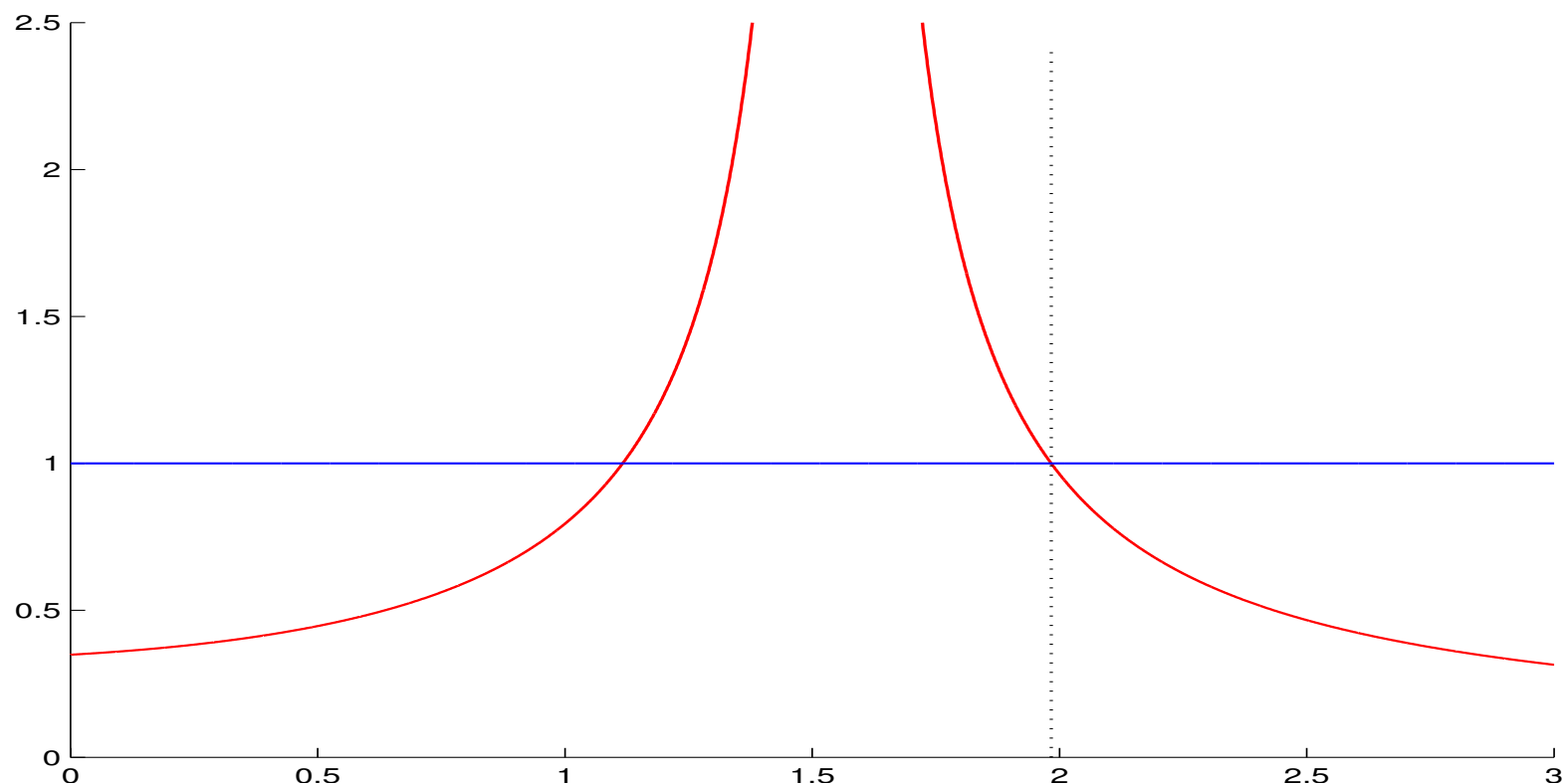
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 4$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

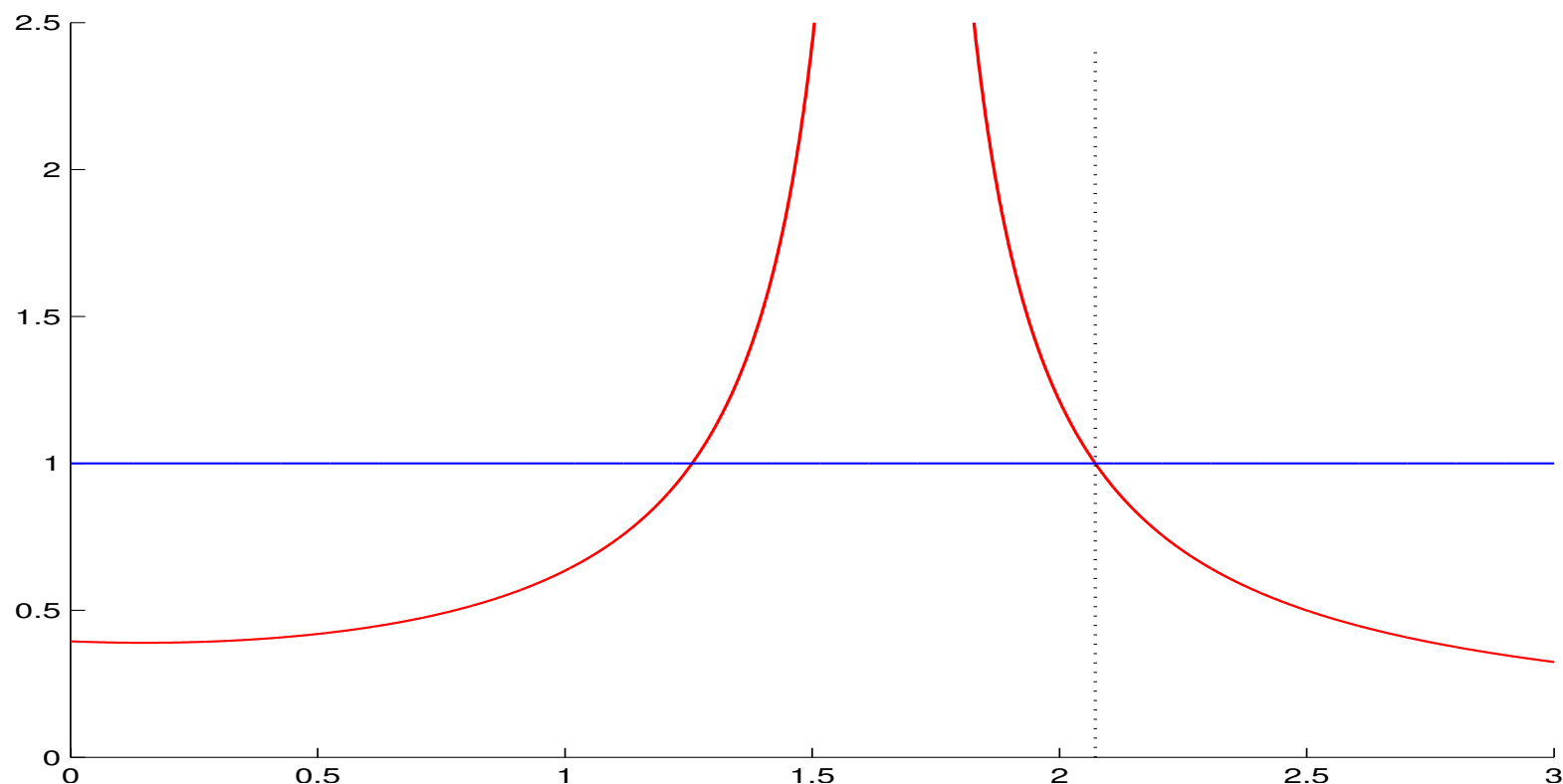
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 5$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

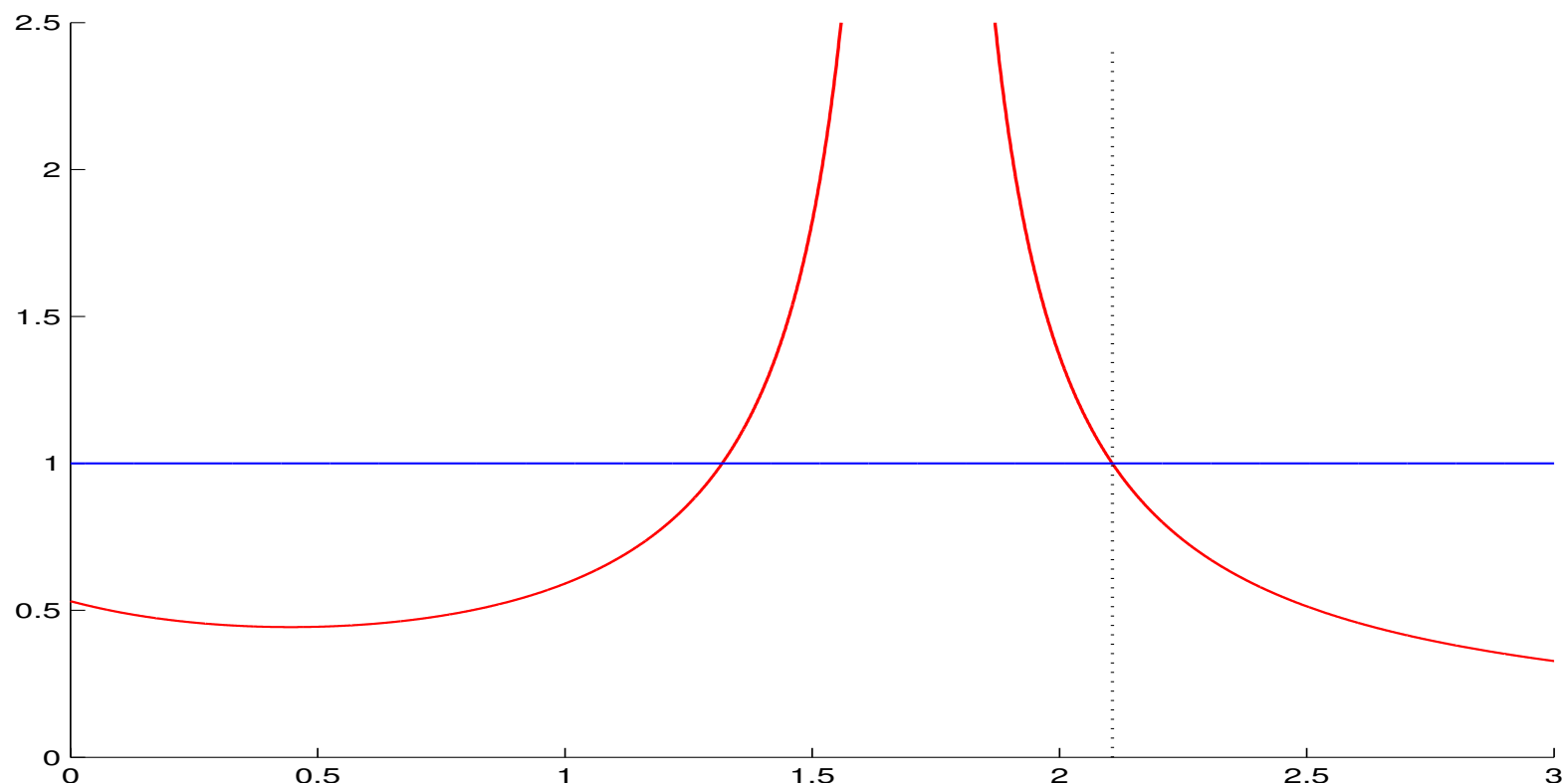
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 6$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

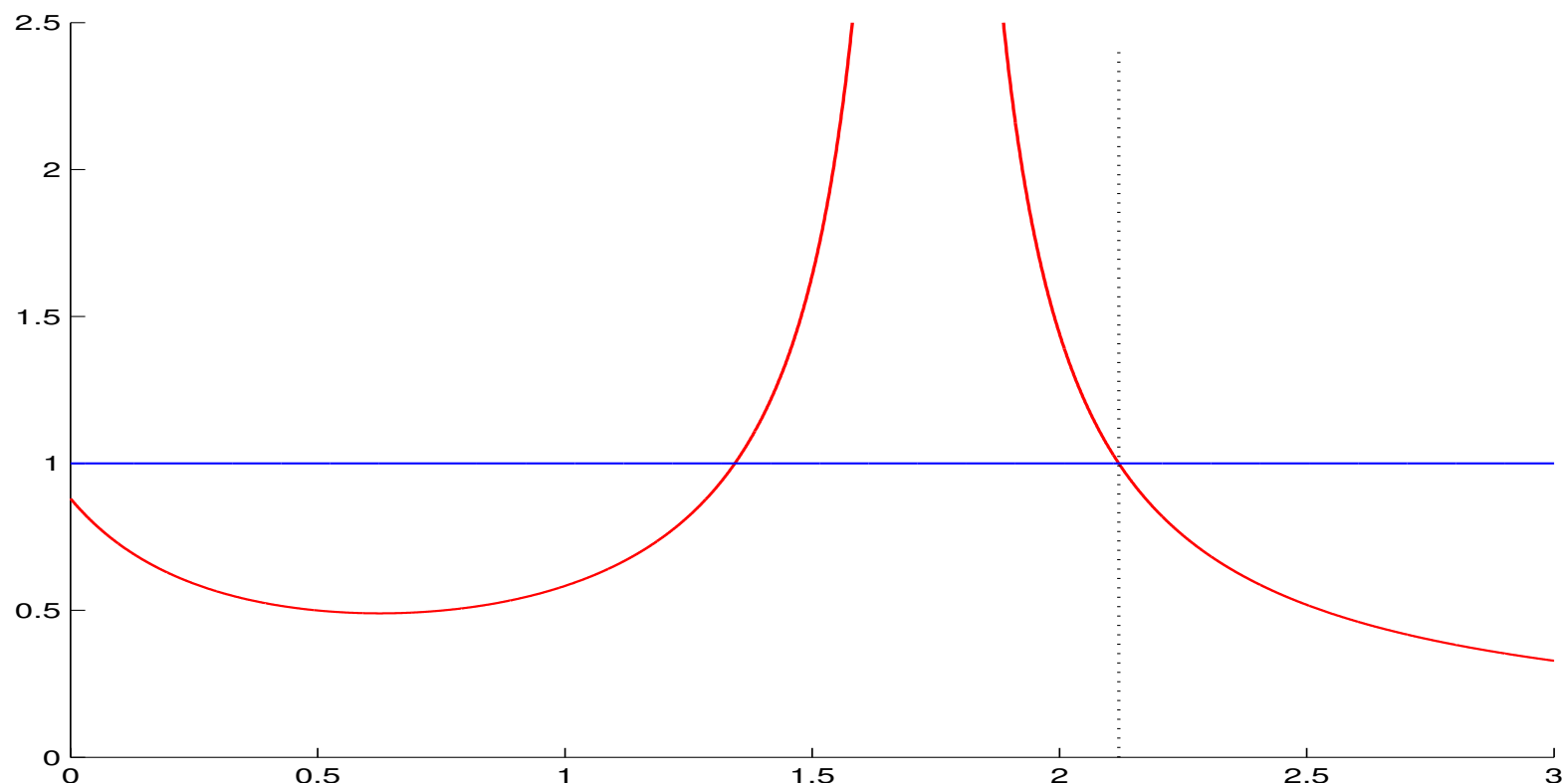
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 7$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

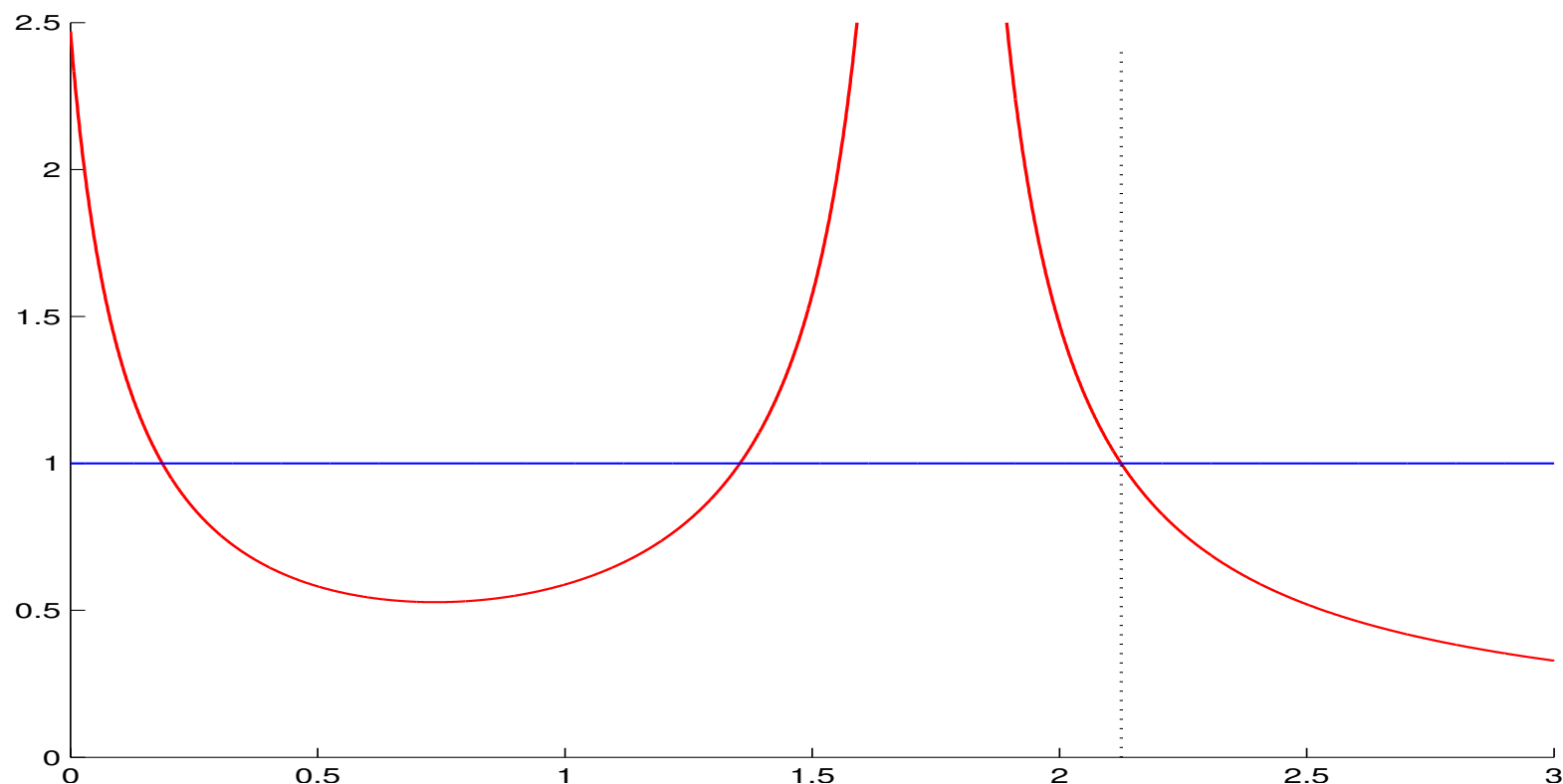
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 8$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

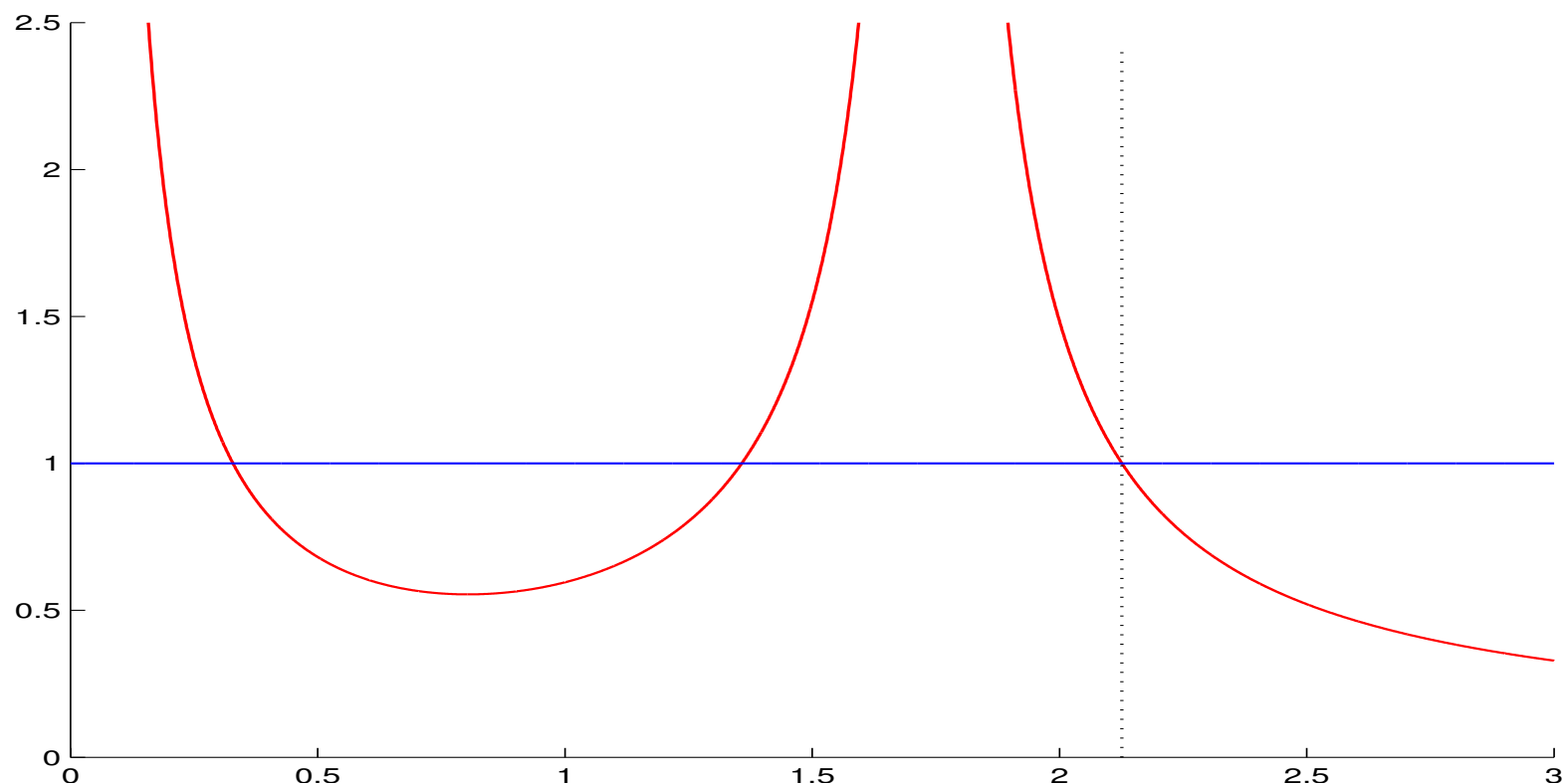
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 9$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

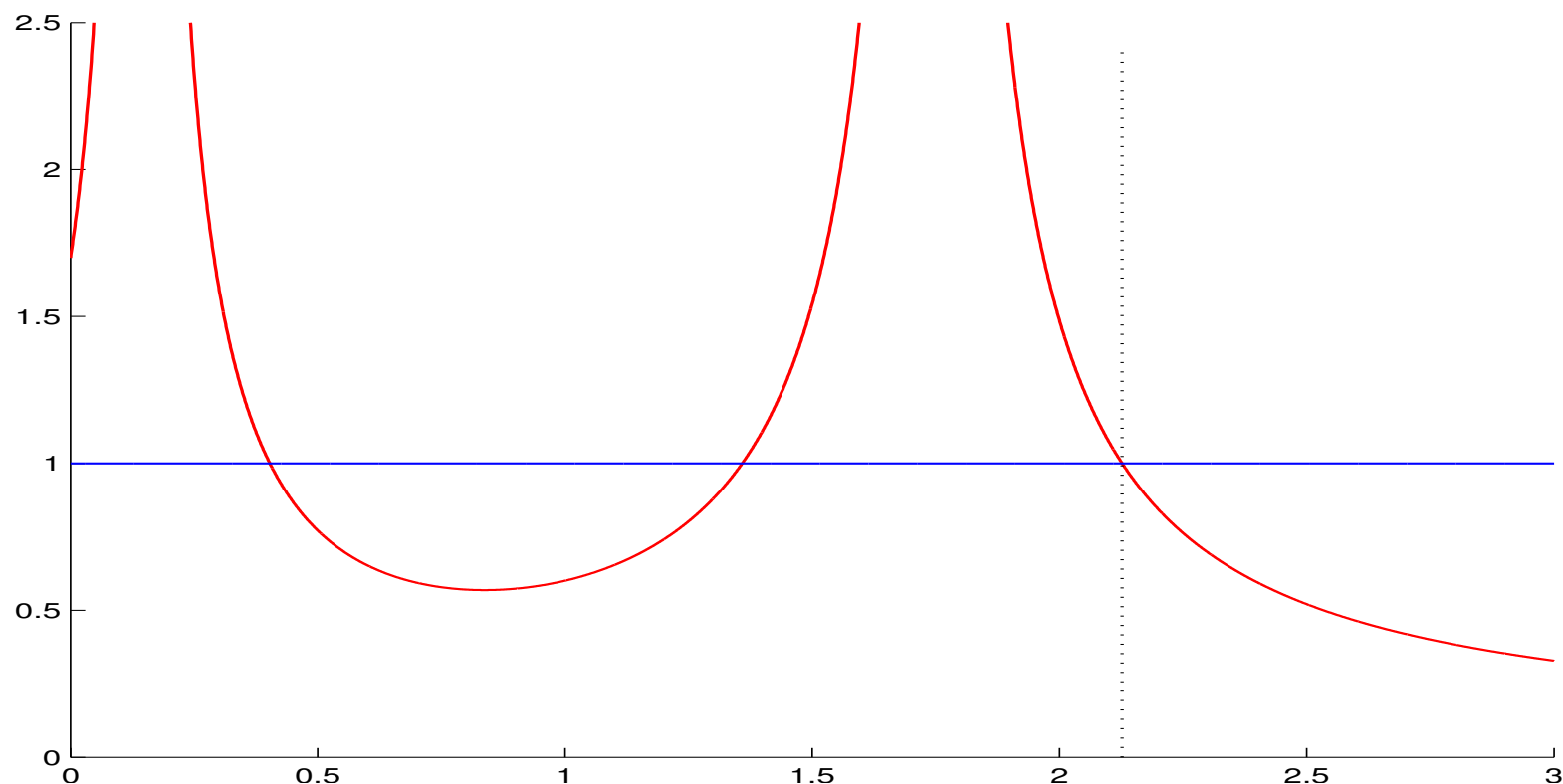
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 10$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

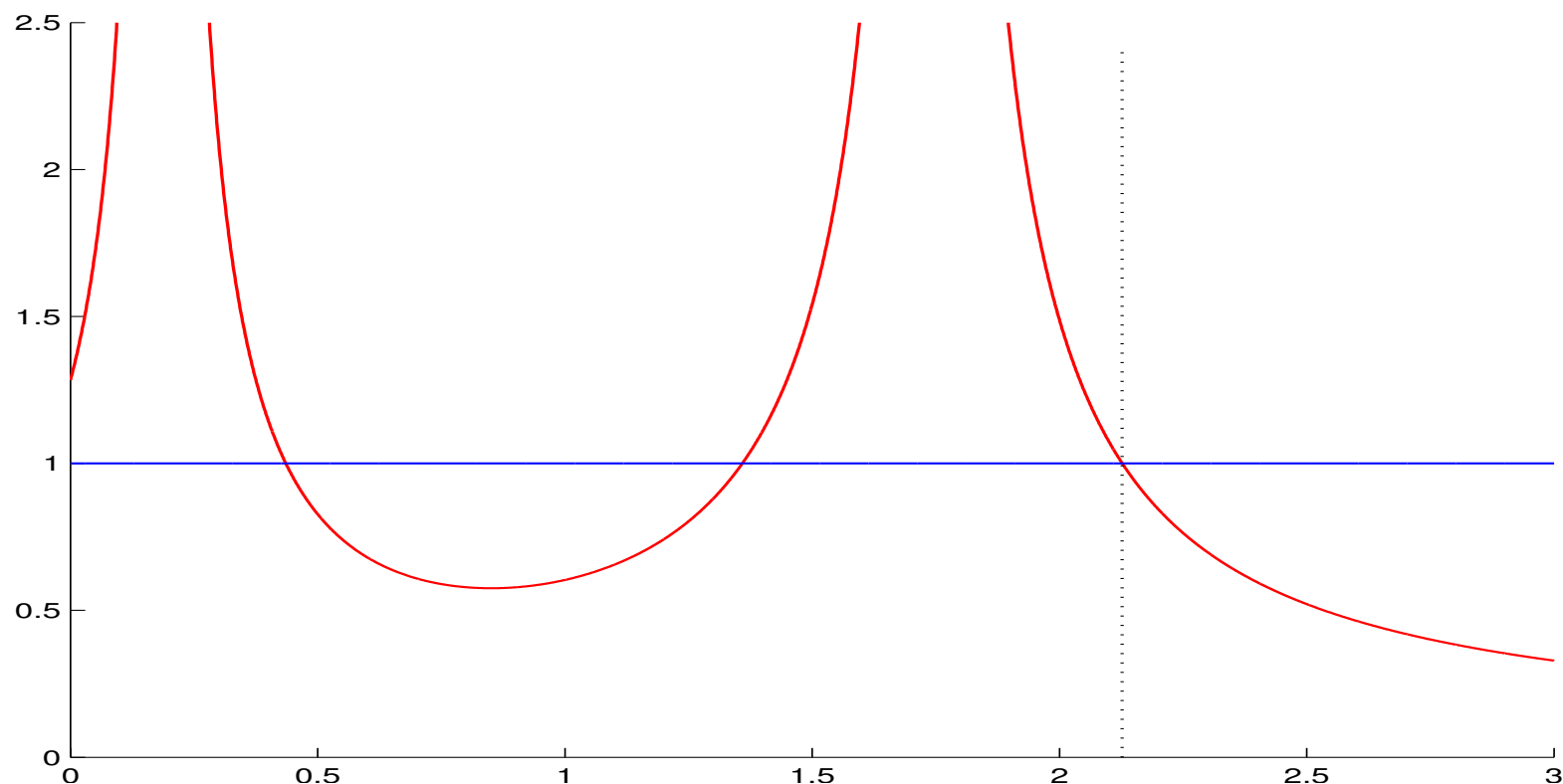
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 11$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

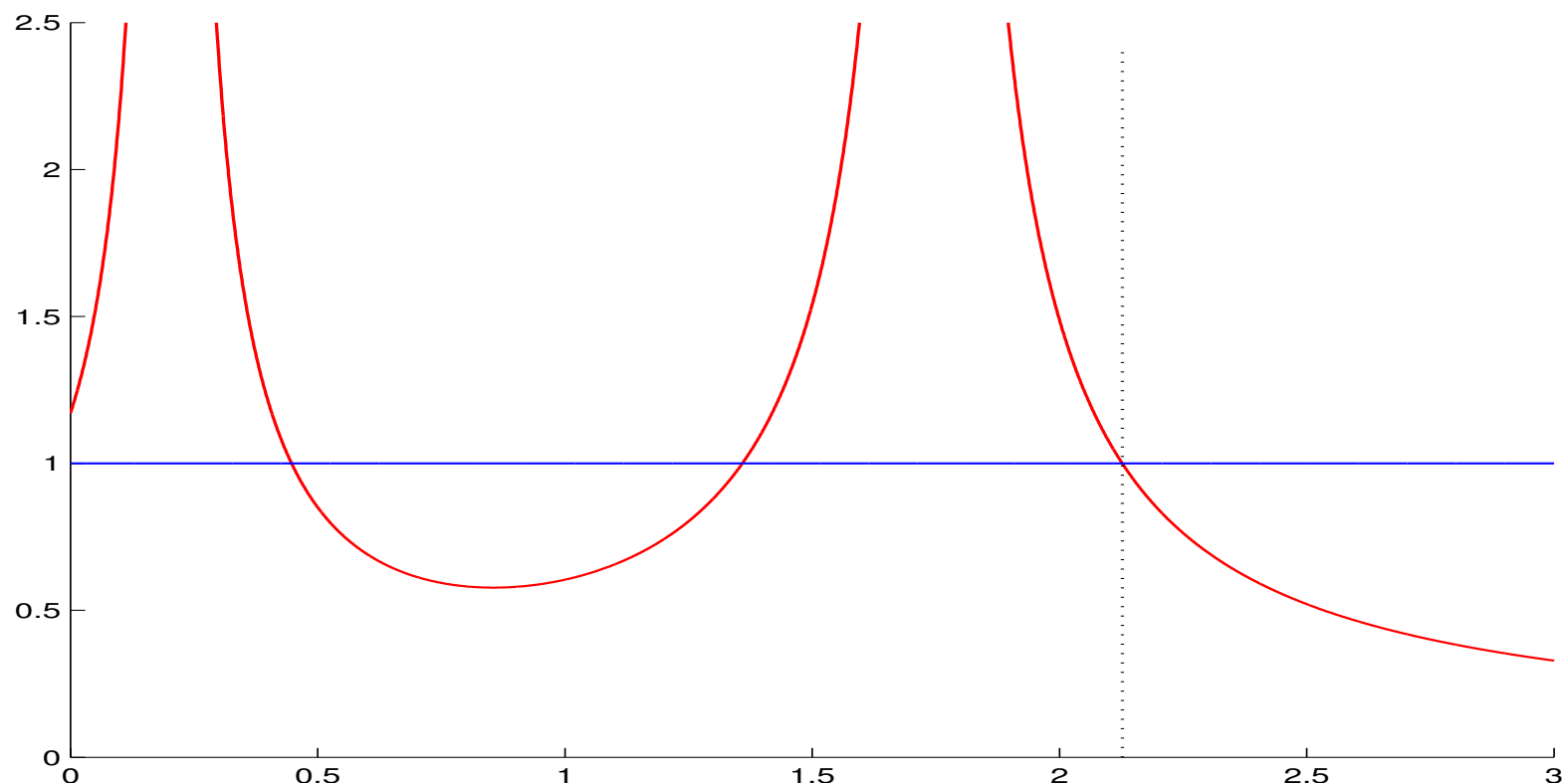
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 12$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

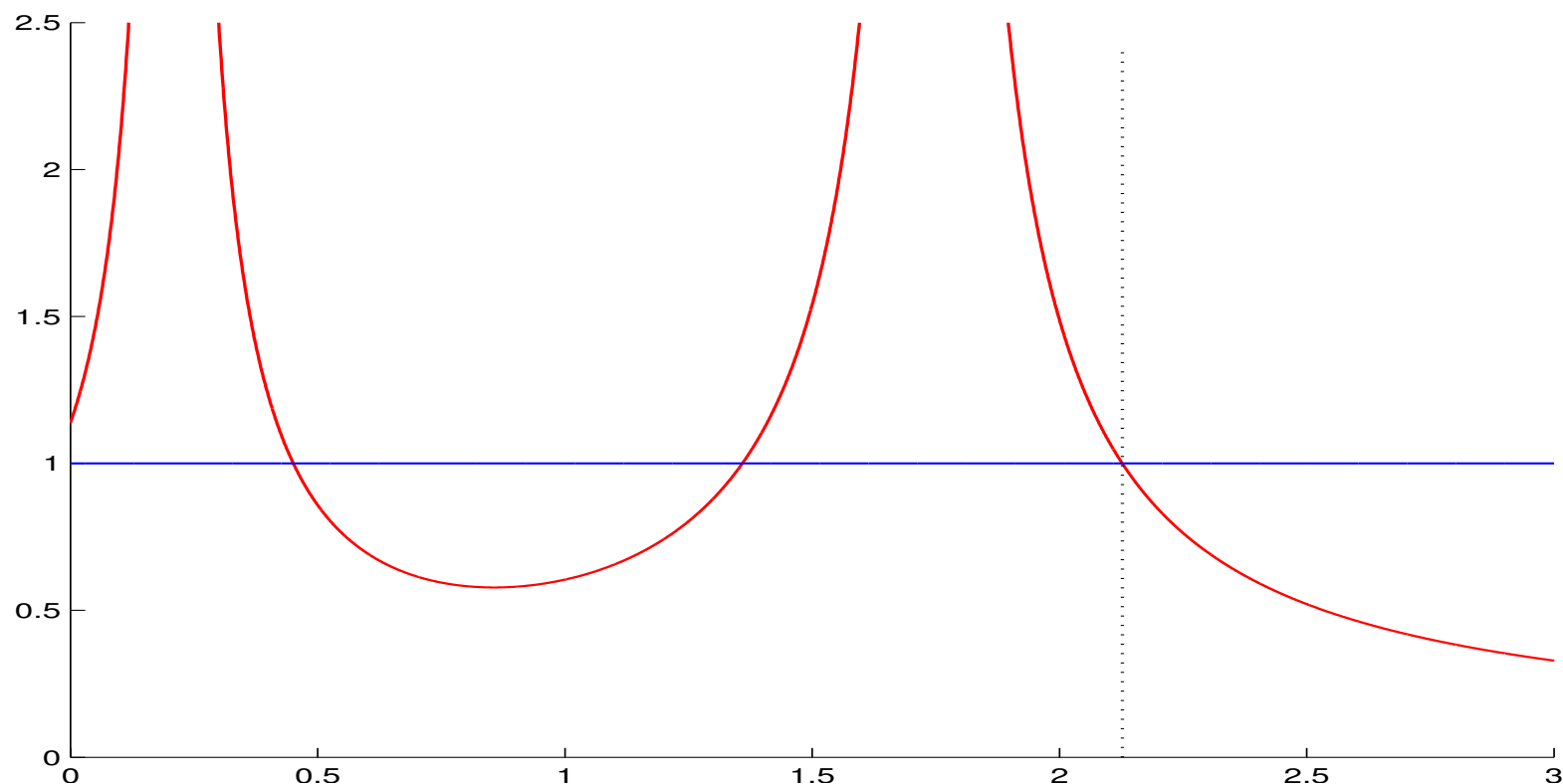
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 13$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

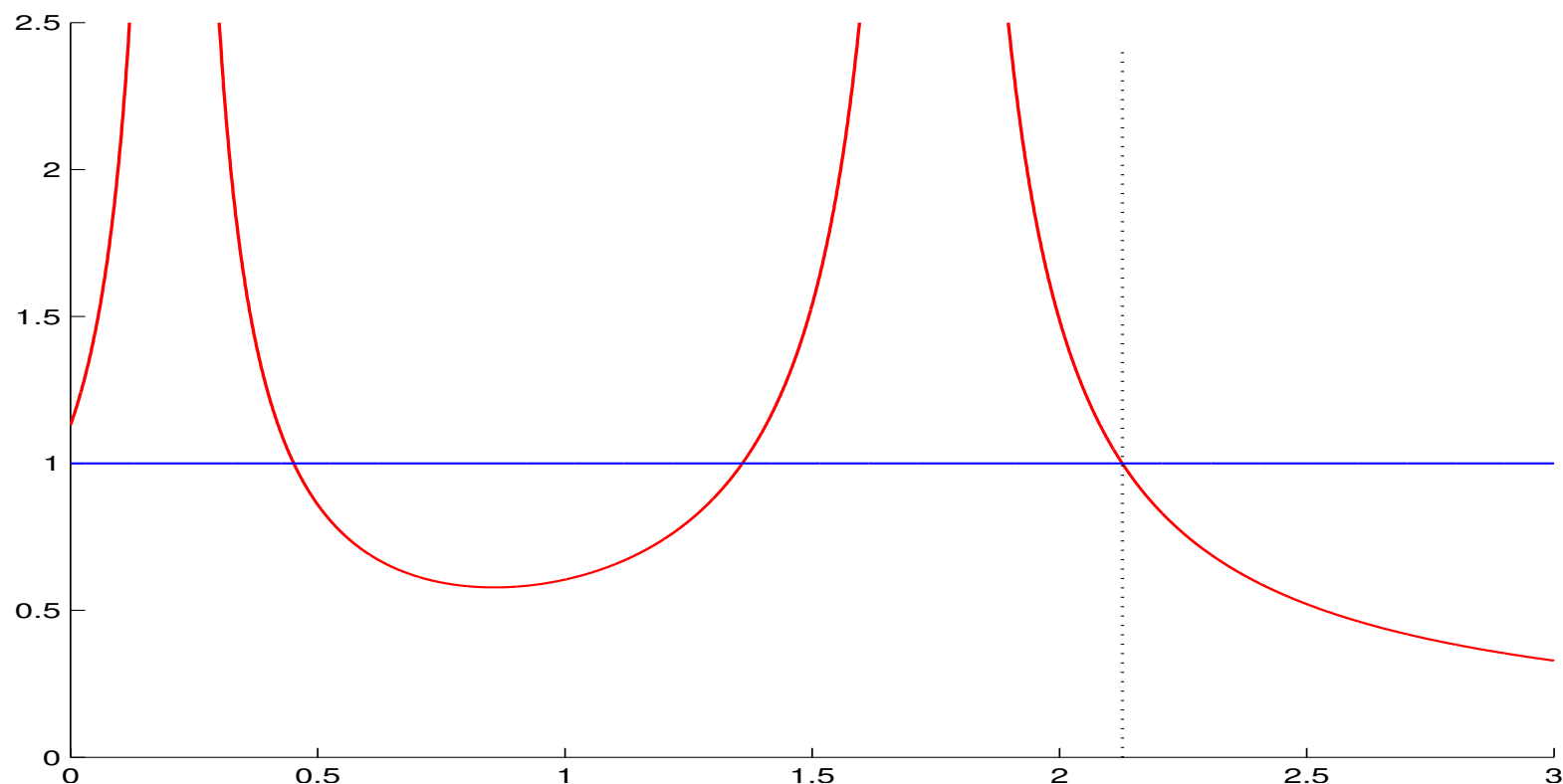
$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$



Iteration $k = 14$ for trust-region problem ($\tau(\lambda) = 1$)

How the secular equation evolves

$$(T_k + \lambda I)x_k(\lambda) = \|b\|_2 e_1 \quad \text{and} \quad \|x_k(\lambda)\|_2 = \tau(\lambda)$$




Iteration $k = 15$ for trust-region problem ($\tau(\lambda) = 1$)

Conclusions

- many important problems involve secular equations
- nonlinear transformation of variables and/or equations very useful
- high-order methods often pay off
- can be solved by either factorization or subspace iteration
- future advances to both large-scale linear equation and eigen solvers will be beneficial

Conclusions

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- nonlinear transformation of variables and/or equations very useful
- high-order methods often pay off
- can be solved by either factorization or subspace iteration
- future advances to both large-scale linear equation and eigen solvers will be beneficial
- freely available optimization software as part of **GALAHAD** 
- TRS/RQS** for direct solution of trust-region/regularization subproblems (G., Robinson, Thorne, 2011)
- GLTR/GLRT** for iterative solution of trust-region/regularization subproblems (G., Lucidi, Roma, Toint, 1999, Cartis, G., Toint, 2011)
- LSTR/LSRT/L2RT** for iterative solution of trust-region/regularization of least-squares subproblems (Cartis, G., Toint, 2010)